

COMPLEXES OF DIRECTED TREES OF COMPLETE MULTIPARTITE GRAPHS

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ABSTRACT. For every directed graph D we consider the complex of all directed subforests $\Delta(D)$. The investigation of these complexes was started by D. Kozlov. We generalize a result of Kozlov and prove that complexes of directed trees of complete multipartite graphs are shellable. We determine the h -vector of $\Delta(\vec{K}_{m,n})$ and the homotopy type of $\Delta(\vec{K}_{n_1, n_2, \dots, n_k})$.

1. Introduction

A *directed tree* is a tree in which one vertex is selected as the *root* and all edges are oriented away from the root. If $T = (V(T), E(T))$ is a directed tree with root r , then for every $x \in V(T)$ there exists a unique directed path from r to x . We say that a vertex y is *below* vertex x in a directed tree T if there exists a unique directed path from x to y . A *directed forest* is a family of disjoint directed trees. In this paper we write \vec{xy} for a directed edge from x to y .

An *abstract simplicial complex* Δ is a collection of finite nonempty sets such that $A \subseteq B \in \Delta \Rightarrow A \in \Delta$. The element A of Δ is called a *face* (*simplex*) of Δ and its dimension is $|A| - 1$. The *vertex set* of Δ is the union of all faces of Δ . The *dimension* of the complex Δ is defined as the largest dimension of any of its faces. A *facet* of Δ is any simplex that is not a face of any larger simplex of Δ . A simplicial complex is *pure* if every of its facets has the same dimension. We denote the number of i -dimensional faces of Δ by f_i , and $f(\Delta) = (f_{-1}, f_0, \dots, f_{d-1}, f_d)$ is called the *f-vector*. A new invariant, the *h-vector* of a d -dimensional complex Δ is $h(\Delta) = (h_0, h_1, \dots, h_d, h_{d+1})$ defined by the formula

$$h_k = \sum_{i=0}^k (-1)^{k-i} \binom{d+1-i}{d+1-k} f_{i-1}.$$

We refer the reader to [8] for definitions of topological concepts used in this paper.

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DEFINITION 1.1. Let D be a directed graph. The vertices of the *complex of directed trees* $\Delta(D)$ are oriented edges of D . The faces of $\Delta(D)$ are all directed forests that are subgraphs of D .

The question of shellability of complexes of directed trees was posed by R. Stanley. Kozlov in [6] showed that the existence of a complete source in a directed graph provides a shelling of its complex of directed trees. The complex of directed trees of a graph G is recognized in [3] as a discrete Morse complex of a 1-dimensional complex. These complexes are also studied in [4] and [7].

Geometrically, a shelling of a cell complex is a way of gluing it together from its maximal cells in a well-behaved way. In this paper we use the following definition of shellability for pure simplicial complexes.

DEFINITION 1.2. A pure simplicial complex Δ is *shellable* if there exists a linear ordering (*shelling order*) F_1, F_2, \dots, F_k of facets of Δ such that for every $i < j \leq k$ there exist some $l < j$ and a vertex v of F_j such that

$$F_i \cap F_j \subseteq F_l \cap F_j = F_j \setminus \{v\}.$$

For a fixed shelling order F_1, F_2, \dots, F_k of Δ , the *restriction* $\mathcal{R}(F_j)$ of the facet F_j is defined by $\mathcal{R}(F_j) = \{v \text{ is a vertex of } F_j : F_j \setminus \{v\} \subset F_i \text{ for some } 1 \leq i < j\}$. The *type* of the facet F in the given shelling order is the cardinality of $\mathcal{R}(F)$, i.e., $\text{type}(F) = |\mathcal{R}(F)|$. If we build up Δ according to a shelling order, then $\mathcal{R}(F)$ is the unique minimal new face that appears when we add the facet F . For a shellable simplicial complex we have the following combinatorial interpretation of its h -vector: $h_k(\Delta) = |\{F \text{ is a facet of } \Delta : \text{type}(F) = k\}|$. Further, we know that a shellable d -dimensional simplicial complex Δ is homotopy equivalent to a wedge of h_d spheres of dimension d . A set of maximal simplices \mathcal{G} of Δ is a *generating set* of simplices if the removal of interiors of all simplices from \mathcal{G} makes Δ contractible. For a shellable simplicial complex Δ the set of simplices $\mathcal{G} = \{F \text{ is a facet of } \Delta : \mathcal{R}(F) = F\}$ is a generating set of Δ , i.e., the simplicial complex $\Delta \setminus (\bigcup_{F \in \mathcal{G}} F)$ is contractible.

For more information on shellability see [1], [2] and chapter 8 of [9].

2. Graphs with a dominant pair

Kozlov in [6] used the following variant of a shelling. Let $\mathcal{F}(\Gamma)$ denote the set of all facets of a pure simplicial complex Γ . Assume that we can partition $\mathcal{F}(\Gamma)$ into the blocks $\mathcal{F}_0, \mathcal{F}_1, \mathcal{F}_2, \dots, \mathcal{F}_m$ such that the following holds:

$$(2.1) \quad \begin{aligned} &|\mathcal{F}_0| = 1; \text{ for all } i \leq j \text{ and two different facets } F \in \mathcal{F}_i, F' \in \mathcal{F}_j, \\ &\text{there exists } k < j, \text{ a facet } F'' \in \mathcal{F}_k, \text{ and a vertex } v \in F' \\ &\text{such that } F \cap F' \subseteq F'' \cap F' = F' \setminus \{v\}. \end{aligned}$$

It is easy to check that any linear order that refines partition $\mathcal{F}_0, \mathcal{F}_1, \mathcal{F}_2, \dots, \mathcal{F}_m$ (for $i < j$ we list facets from \mathcal{F}_i before facets from \mathcal{F}_j) is a shelling order of Γ in the sense of Definition 1.2.

A vertex c is a *complete source* in a digraph D if $\vec{cx} \in E(D)$ for all $x \in V(D) \setminus \{c\}$. The partition of $\mathcal{F}(\Delta(D))$ defined by Kozlov in [6] substantially uses

the out-degree $d_T(c) = |\{x : \overrightarrow{cx} \in E(T)\}|$ of c :

$$\mathcal{F}_i = \{T \in \mathcal{F}(\Delta(D)) : d_T(c) - 1 - i\}.$$

It is not complicated to prove for a digraph with a complete source, that the above partition of the facets of $\Delta(D)$ satisfies the condition described in (2.1) (see the proof of Theorem 3.1 in [6]).

We now describe a broader family of graphs whose complexes of directed trees are shellable. For a directed graph D and a vertex $u \in V(D)$ we set $N(u) = \{x \in V(D) : \overrightarrow{ux} \in E(D)\}$. We say that a directed graph D has a *dominant pair of vertices* if there exist vertices $u, v \in V(D)$ such that

- (i) $V(D) = N(u) \cup N(v)$. Therefore, we have that $\overrightarrow{uv}, \overrightarrow{vu} \in E(D)$.
- (ii) For all $x \in N(u) \setminus N(v)$ and $y \in N(v) \setminus N(u)$ we have that $\overrightarrow{xy}, \overrightarrow{yx} \in E(D)$.

THEOREM 2.1. *If a directed graph D has a dominant pair of vertices the complex $\Delta(D)$ is shellable.*

PROOF. We will define a partition of the facets of $\Delta(D)$ and show that this partition satisfies (2.1). Recall that facets of the complex $\Delta(D)$ correspond to subtrees of D .

Let D be a graph with a dominant pair of vertices u, v . For a directed tree T with the root r let $h_T(x)$ denotes the length of the unique directed path from r to x . We classify directed trees of D by using $d_T(u)$, $d_T(v)$, $h_T(u)$ and $h_T(v)$. The trees of D in which the above defined parameters are the same form a block

$$\mathcal{F}_{p,q,r,s} = \{T : d_T(u) = p, d_T(v) = q, h_T(u) = r, h_T(v) = s\}$$

of our partition of the facets of $\Delta(D)$. We say that $\mathcal{F}_{p,q,r,s}$ is before $\mathcal{F}_{p',q',r',s'}$, and write $\mathcal{F}_{p,q,r,s} < \mathcal{F}_{p',q',r',s'}$, if and only if $p > p'$, or $p = p'$ and $q > q'$, or $p' = p$, $q = q'$ and $r < r'$, or $p' = p$, $q = q'$, $r = r'$ and $s < s'$. Note that the first block in this partition $\mathcal{F}_{|N(u)|, |N(v) \setminus N(u)| - 1, 0, 1}$ contains only the tree with the edge set $\{\overrightarrow{ux} : x \in N(u)\} \cup \{\overrightarrow{vy} : y \in N(v) \setminus (N(u) \cup \{u\})\}$.

Now, we consider two different directed trees $T \in \mathcal{F}_{p,q,r,s}$, $T' \in \mathcal{F}_{p',q',r',s'}$ such that $\mathcal{F}_{p,q,r,s} \leq \mathcal{F}_{p',q',r',s'}$. Assume that the edges $E(T) \cap E(T')$ span a directed forest $F = T_1 \cup T_2 \cup \dots \cup T_m$, and let r_i denote the root of T_i . Assume that r_{i_0} is the root of T' . Note that $E(T') \setminus E(T)$ contains $m - 1$ edges of the form $\overrightarrow{xr_i}$, where $i \neq i_0$.

The following analysis will show that there is a tree T'' such that T, T' and T'' satisfy the conditions described in (2.1). First, we consider the case when u, v are in the same tree of F (w.l.o.g. we assume $v, u \in T_1$).

- (1) If the root of T' is a vertex $r_i \neq r_1$, then there exists $\overrightarrow{xr_1} \in E(T') \setminus E(T)$ such that $x \neq u$ and $x \neq v$. We set $T'' = T' \setminus \{\overrightarrow{xr_1}\} \cup \{\overrightarrow{ur_1}\}$ (if $r_i \in N(u)$) or $T'' = T' \setminus \{\overrightarrow{xr_1}\} \cup \{\overrightarrow{vr_1}\}$ (if $r_i \notin N(u)$).
- (2) Assume that r_1 is the root of T' . If there is a vertex $r_j \in N(u)$ for some $j > 1$, the assumption $d_T(u) \geq d_{T'}(u)$ guarantees that there exists an edge $\overrightarrow{xr_j} \in E(T') \setminus E(T)$ such that $x \neq u$, $j > 1$ and $r_j \in N(u)$. Otherwise, if all of the edges $\overrightarrow{ur_i}$ (for all $r_i \in N(u)$, $i > 1$) are contained in $E(T')$, then we obtain that $d_T(u) < d_{T'}(u)$. In the above described situation we set

$$T'' = T' \setminus \{\overrightarrow{xr_i}\} \cup \{\overrightarrow{ur_i}\}.$$

If $r_i \in N(v) \setminus N(u)$ for all $i = 2, 3, \dots, m$, then there exists $\overrightarrow{yr_i} \in E(T') \setminus E(T)$, such that $y \neq v$ (otherwise we obtain that $d_T(u) = d_{T'}(u)$ and $d_T(v) < d_{T'}(v)$). Then we set $T'' = T' \setminus \{\overrightarrow{yr_i}\} \cup \{\overrightarrow{vr_i}\}$.

Now, we consider the situation when the vertices u and v belong to different trees of F (w.l.o.g. we assume that $u \in T_1$ and $v \in T_2$).

- (3) If the root of T' is r_i , $r_i \neq r_1$ and $r_i \in N(u)$, then there exists $\overrightarrow{xr_1} \in E(T') \setminus E(T)$ such that $x \neq u$ and we set $T'' = T' \setminus \{\overrightarrow{xr_1}\} \cup \{\overrightarrow{ur_1}\}$.
- (4) If the root r_i of T' (again $r_i \neq r_1$) is not contained in $N(u)$ and $r_i \neq r_2$, then there exists $\overrightarrow{xr_2} \in E(T') \setminus E(T)$. If $x \neq u$ we set $T'' = T' \setminus \{\overrightarrow{xr_2}\} \cup \{\overrightarrow{vr_2}\}$.
If $x = u$, then v is below u and there exists $\overrightarrow{yr_1} \in E(T')$ such that $y \neq u$ and $y \neq v$. In that case we set $T'' = T' \setminus \{\overrightarrow{yr_1}\} \cup \{\overrightarrow{vr_1}\}$.
- (5) If r_2 is the root of T' (recall that $r_2 \in N(v) \setminus N(u)$ and therefore $r_2 \neq v$), then there exists an edge $\overrightarrow{xr_1} \in E(T') \setminus E(T)$. If $r_1 \in N(v)$ and $x \neq v$, we set $T'' = T' \setminus \{\overrightarrow{xr_1}\} \cup \{\overrightarrow{vr_1}\}$.

If $x = v$ (and therefore $r_1 \in N(v)$), then we find an edge $\overrightarrow{yr_i} \in E(T') \setminus E(T)$ such that $y \neq u$, $i > 2$ and $r_i \in N(u)$ (or $\overrightarrow{zr_j} \in E(T') \setminus E(T)$ such that $z \neq v$, $r_j \in N(v) \setminus N(u)$, $j > 2$) by using the same arguments as in the proof of (2). Then we set $T'' = T' \setminus \{\overrightarrow{yr_i}\} \cup \{\overrightarrow{ur_i}\}$ or $T'' = T' \setminus \{\overrightarrow{zr_j}\} \cup \{\overrightarrow{vr_j}\}$.

If $r_1 \notin N(v)$, then there exists $\overrightarrow{xr_1} \in E(T') \setminus E(T)$, $x \neq u$, $x \neq v$ and we set $T'' = T' \setminus \{\overrightarrow{xr_1}\} \cup \{\overrightarrow{r_1r_2}\}$. In that case we obtain that $d_{T''}(u) = d_{T'}(u)$, $d_{T''}(v) = d_{T'}(v)$ and $h_{T''}(u) < h_{T'}(u)$.

If the root of T' is r_1 , then we have the following possibilities.

- (6) There exists $i > 1$ such that $r_i \in N(u)$. Because we have that $d_T(u) \geq d_{T'}(u)$ it follows that there exists $\overrightarrow{xr_j} \in E(T') \setminus E(T)$ such that $x \neq u$, $r_j \in N(u)$ and $j > 1$. In that case we set $T'' = T' \setminus \{\overrightarrow{xr_j}\} \cup \{\overrightarrow{ur_j}\}$.
- (7) If all vertices r_i for $i = 2, 3, \dots, m$ are contained in $N(v) \setminus N(u)$ and $r_1 \in N(v)$ (recall that r_1 is the root of T'), then there exists $\overrightarrow{xr_2} \in E(T') \setminus E(T)$, $x \neq u$ and we set $T'' = T' \setminus \{\overrightarrow{xr_2}\} \cup \{\overrightarrow{vr_1}\}$.
- (8) Finally, we assume that $r_1 \in N(u) \setminus N(v)$, $r_i \in N(v) \setminus N(u)$ for all $i > 1$, and r_1 is the root of T' . In that case we have that $d_T(u) = d_{T'}(u)$.

If $m > 2$, from $d_T(v) \geq d_{T'}(v)$ we conclude that there exists $r_i \in N(v)$ for $i > 2$ and an edge $\overrightarrow{yr_i} \in E(T') \setminus E(T)$ such that $y \neq v$. Then we set $T'' = T' \setminus \{\overrightarrow{yr_i}\} \cup \{\overrightarrow{vr_i}\}$.

For $m = 2$, we again consider the edge $\overrightarrow{xr_2} \in E(T') \setminus E(T)$. If $x = r_1$, we have that $d_T(u) = d_{T'}(u)$, $d_T(v) = d_{T'}(v)$ and

$$T = T' \setminus \{\overrightarrow{r_1r_2}\} \cup \{\overrightarrow{zr_1}\} \text{ or } T = T' \setminus \{\overrightarrow{r_1r_2}\} \cup \{\overrightarrow{yr_2}\}.$$

But, then we have that $h_T(u) > h_{T'}(u)$ or $h_T(u) = h_{T'}(u)$, $h_T(v) > h_{T'}(v)$ which is a contradiction with the assumption. Therefore, in this case ($m = 2$) we have that $x \neq r_1$. If we set $T'' = T' \setminus \{\overrightarrow{xr_2}\} \cup \{\overrightarrow{r_1r_2}\}$, then we obtain that $d_{T''}(u) = d_{T'}(u)$, $d_{T''}(v) = d_{T'}(v)$, $h_{T''}(u) = h_{T'}(u)$ and $h_{T''}(v) < h_{T'}(v)$. \square

3. A complete multipartite graph

Let K_{n_1, n_2, \dots, n_k} denote a complete multipartite graph. Assume that its vertex set is $V = V_1 \cup V_2 \cup \dots \cup V_k$, where $|V_i| = n_i$. Furthermore, we assume that all sets V_i are linearly ordered. We may choose one vertex in V_1 and one in V_2 and denote them by 1 and -1 .

Let $\vec{K}_{n_1, n_2, \dots, n_k}$ denote a directed graph obtained from a complete multipartite graph K_{n_1, n_2, \dots, n_k} when one replaces all edges by pairs of directed edges going in opposite directions. Note that 1, -1 is a dominant pair of vertices in $\vec{K}_{n_1, n_2, \dots, n_k}$ and from Theorem 2.1 we know that $\Delta(\vec{K}_{n_1, n_2, \dots, n_k})$ is shellable. We use a slight modification of the algorithm described in [5] to encode directed trees in $\vec{K}_{n_1, n_2, \dots, n_k}$.

REMARK 3.1. For each directed tree T of $\vec{K}_{n_1, n_2, \dots, n_k}$ we associate the set of sequences $\{C_0, C_1, \dots, C_k\}$ of the vertex set such that

- (i) The length of the sequence C_0 is $k - 1$ and any $x \in V$ can occur in C_0 .
- (ii) For any $i > 0$ the length of C_i is $n_i - 1$ and C_i contains vertices from $V \setminus V_i$.

Let r denote the root of T . For a vertex $v \in V$, $v \neq r$, let $U_T(v)$ denote the unique vertex u such that $\vec{uv} \in E(T)$. We say that the *depth* of a vertex v in T (denoted by $\text{depth}(v)$) is the length of the longest directed path from v to a leaf of T . For all $i = 1, 2, \dots, k$ let v'_i denote the vertex from V_i with the maximal depth in T (if there are more than one vertex in V_i with maximal depth for v'_i , we choose the greatest one among them in the linear order of V_i).

If the root of T is a vertex that belongs to V_{i_0} , then we have that $v'_{i_0} = r$.

The sequence C_0 contains vertices $\{U_T(v'_i) : i \neq i_0\}$, and the vertex $U_T(v'_j)$ is before $U_T(v'_s)$ in C_0 if and only if $\text{depth}(v'_j) < \text{depth}(v'_s)$ or $\text{depth}(v'_j) = \text{depth}(v'_s)$ and $j < s$. For any $i > 0$ the entries of the sequence C_i are $n_i - 1$ vertices $\{U_T(v) : v \in V_i, v \neq v'_i\}$ and we order the set of these vertices in the same way as in C_0 . Vertices from V_j that appear in C_i and have the same depth, we order in C_i by using the linear order defined on V_j . We say that $\{C_0, C_1, \dots, C_k\}$ is the code for the tree T . The proof that the map $T \mapsto \{C_0, C_1, \dots, C_k\}$ is a bijection, as well as more details about this construction can be found in [5].

It is easily seen from the above remark that there are

$$n^{k-1}(n - n_1)^{n_1-1}(n - n_2)^{n_2-1} \dots (n - n_k)^{n_k-1}$$

directed trees in $\vec{K}_{n_1, n_2, \dots, n_k}$. These are the facets of $\Delta(\vec{K}_{n_1, n_2, \dots, n_k})$.

THEOREM 3.1. *The h -vector of $\Delta(\vec{K}_{m, n})$ is given by*

$$\begin{aligned} h_k(\Delta(K_{m, n})) &= \sum_{p+q=k} \binom{m-1}{p} (n-1)^p \binom{n-1}{q} (m-1)^q \\ &\quad + (m+n-1) \sum_{p+q=k-1} \binom{m-1}{p} (n-1)^p \binom{n-1}{q} (m-1)^q. \end{aligned}$$

PROOF. Note that $\Delta(\vec{K}_{m,n})$ is $(m+n-2)$ -dimensional complex. We consider the shelling order of $\Delta(\vec{K}_{m,n})$ described in Theorem 2.1. Recall that

$$\mathcal{R}(T) = \{\vec{xy} \in E(T) : E(T) \setminus \{\vec{xy}\} \subset E(S) \text{ for some tree } S \text{ that precedes } T\}.$$

In other words, an edge $\vec{xy} \in E(T)$ is in $\mathcal{R}(T)$ if it can be replaced with another edge $\vec{zw} \notin E(T)$ such that $(T \setminus \{\vec{xy}\}) \cup \{\vec{zw}\}$ is a new directed tree which precedes T in considered shelling order.

It is easy to check that the following statements hold:

- (i) The restriction $\mathcal{R}(T)$ does not contain the edge $\vec{1x}$. A replacement of $\vec{1x}$ will decrease the out-degree of 1.
- (ii) A replacement of the edge $\vec{-1x}$ in T will decrease the out-degree of -1 . A new tree $T' = (T \setminus \{\vec{-1x}\}) \cup \{\vec{yz}\}$ precedes T in the considered shelling order only if we increase the out-degree of 1. We can do this if and only if the vertex $y = 1$ is below -1 in T and $z \in V_2$ is the root of T . Other edges $\vec{-1x'}$ can not be replaced.
- (iii) Let r be the root of T . For a vertex $x \in V_1, x \neq 1$, and an edge \vec{xy} we have:
 - (a) if 1 is not below y the tree $(T \setminus \{\vec{xy}\}) \cup \{\vec{1y}\}$ precedes T .
 - (b) If 1 is below y and if r belongs V_2 , we have that $(T \setminus \{\vec{xy}\}) \cup \{\vec{1y}\}$ is before T .
 - (c) If $r \in V_1$ (recall that 1 is below y) we set $S = (T \setminus \{\vec{xy}\}) \cup \{\vec{yr}\}$. Then we have $d_T(1) = d_S(1)$, $d_T(-1) = d_S(-1)$, $h_T(1) > h_S(1)$ and therefore the tree S precedes T .

So, any of the considered edges \vec{xy} is contained in $\mathcal{R}(T)$.

A similar analysis shows that an edge \vec{xy} , where $x \in V_2, x \neq -1$, is contained in $\mathcal{R}(T)$ except when x is the root of T , -1 is below y and 1 is not below y .

From the above remarks we have that for a directed tree T

$$(3.1) \quad \text{type}(T) = m + n - 1 - d_T(1) - d_T(-1)$$

except for the following trees:

- (A1) Trees in which the root r belongs to V_2 and the vertex 1 is below of -1 . The type of a such tree T is $\text{type}(T) = m + n - d_T(1) - d_T(-1)$.
- (A2) Trees in which the root $r \in V_2, r \neq -1$, there exists an edge $\vec{rx} \in E(T)$ such that -1 is below x and 1 is not below x . The type of this tree is $\text{type}(T) = m + n - d_T(1) - d_T(-1) - 2$.

Now, we count the number of trees in $\Delta(\vec{K}_{m,n})$ with given $d_T(1) + d_T(-1)$. Let $\{r\}, C_1, C_2$ be the set of sequences of vertices associated to a tree T (r is the root of T) in Remark 3.1. We set

$$p = |\{x \in C_1 : x \neq -1\}|, \quad q = |\{y \in C_2 : y \neq 1\}|.$$

From Remark 3.1 we obtain that there are

$$(m+n-2) \sum_{p+q=k-1} \binom{m-1}{p} (n-1)^p \binom{n-1}{q} (m-1)^q \\ + 2 \sum_{p+q=k} \binom{m-1}{p} (n-1)^p \binom{n-1}{q} (m-1)^q$$

directed trees in $\vec{K}_{m,n}$ such that $d_T(1) + d_T(-1) = m+n-1-k$. Note that the summands in the second row correspond with the trees in which the root is 1 or -1 . From the relation (3.1) we have that all of these trees are of the type k , except the trees described in (A1) and (A2).

The remaining trees of $\vec{K}_{m,n}$ of the type k are all

- (B1) trees described in (A2) in which $d_T(1) + d_T(-1) = m+n-2-k$; or
- (B2) trees described in (A1) in which $d_T(1) + d_T(-1) = m+n-k$.

Let T be a directed tree as considered in (B1). If $\vec{r\hat{y}}$ is the first edge on the path from r to 1, then $T' = (T \setminus \{\vec{r\hat{y}}\}) \cup \{-1\hat{y}\}$ is a tree as described in (A1). Note that the map $T \mapsto T'$ is an injection, and all trees described in (A1) are contained in the image of this map except the trees whose root is -1 . From Remark 3.1 it follows that there are

$$\sum_{p+q=k} \binom{m-1}{p} (n-1)^p \binom{n-1}{q} (m-1)^q$$

trees with -1 as the root and $d_T(1) + d_T(-1) = m+n-1-k$, which should be subtracted while calculating $h_k(\Delta(\vec{K}_{m,n}))$.

Further, if T is a tree described in (A2), and $\vec{r\hat{x}}$ is the first edge of the path from r to 1, then $T' = (T \setminus \{\vec{r\hat{x}}\}) \cup \{-1\hat{x}\}$ is a tree as in (B2). This map is an injection, and a tree from (B2) is not in the image of this map if and only if its root is -1 .

There are

$$\sum_{p+q=k-1} \binom{m-1}{p} (n-1)^p \binom{n-1}{q} (m-1)^q$$

trees with -1 as the root and $d(1) + d(-1) = m+n-k$ that should be added when determining $h_k(\Delta(\vec{K}_{m,n}))$. \square

From the above theorem we obtain that the generating facets for $\Delta(\vec{K}_{m,n})$ are:

- (i) All directed trees of $\vec{K}_{m,n}$ in which the vertices 1 and -1 are leaves and the root of such a tree is a vertex contained in V_1 .
- (ii) All directed trees of $\vec{K}_{m,n}$ in which the root is from V_2 , the vertex 1 is a leaf below -1 , and the out-degree of the vertex -1 in such a tree is one.

COROLLARY 3.1. *The complex $\Delta(\vec{K}_{m,n})$ is homotopy equivalent to a wedge of $(m+n-1)(m-1)^{n-1}(n-1)^{m-1}$ spheres of dimension $m+n-2$.*

THEOREM 3.2. *The complex $\Delta(\vec{K}_{n_1, n_2, \dots, n_k})$ is homotopy equivalent to a wedge of $(n-1)^{k-1}(n-n_1-1)^{n_1-1}(n-n_2-1)^{n_2-1} \dots (n-n_k-1)^{n_k-1}$ spheres of dimension $n-2$.*

PROOF. We use a shelling of $\Delta(\vec{K}_{n_1, n_2, \dots, n_k})$ described in Theorem 2.1 to recognize generating faces. These are

- (A) directed trees in which the vertex 1 is a leaf, and there does not exist an edge $\overrightarrow{-1v}$, for a vertex $v \in V_1$

except the trees of the above form in which

- (A₁) the root is a vertex $v_2 \in V_2$, there is an edge $\overrightarrow{v_2v_1}$ for a vertex $v_1 \in V_1$, the vertex -1 is below v_1 , and 1 is not below v_1 ; and
 (A₂) the root is a vertex $v_1 \in V_1$ and the leaf 1 is below -1 .

Generating facets of $\Delta(\vec{K}_{n_1, n_2, \dots, n_k})$ are also:

- (B) directed trees in which the root is a vertex $r \in V \setminus V_1$, there is only one edge of the form $-1v_1$ for a vertex $v_1 \in V_1$ and 1 is a leaf below v_1 .

Now, we define a map between a subset of the trees of the type (B) and directed trees of type (A₁) or (A₂). If T is a tree of the type (B) with the root $r \in V \setminus V_1$ and $-1 \rightarrow v_1 \rightarrow x \rightarrow y \rightarrow \dots \rightarrow 1$ is the unique path from -1 to 1, then

$$T' = T \setminus \{\overrightarrow{-1v_1}, \overrightarrow{v_1x}\} \cup \{\overrightarrow{xv_1}, \overrightarrow{v_1r}\} \text{ is a tree of type } A_1 \text{ if } x \in V_2,$$

$$T'' = T \setminus \{\overrightarrow{-1v_1}, \overrightarrow{v_1x}\} \cup \{\overrightarrow{-1x}, \overrightarrow{v_1r}\} \text{ is a tree of type } A_2 \text{ if } x \in V \setminus (V_1 \cup V_2).$$

The above map is a bijection that exhausts all trees of type (B) except the trees in which $\overrightarrow{-11}$ is an edge. Therefore, in order to estimate the number of the generating simplices of $\Delta(\vec{K}_{n_1, n_2, \dots, n_k})$ we have to count directed trees in $\vec{K}_{n_1, n_2, \dots, n_k}$ in which

- (*) 1 is a leaf, there are no other edges of the form $\overrightarrow{-1v}$, for a vertex $v \in V_1$;
 or
 (**) 1 is a leaf, $\overrightarrow{-11}$ is an edge, there are no other edges of the form $\overrightarrow{-1v}$, for a vertex $v \in V_1$, and the root is a vertex $r \in V \setminus V_1$.

From Remark 3.1 we obtain that the code of a tree described in (*) or (**) does not contain label -1 in the sequences C_0 at the place reserved for the deepest vertex of V_1 . Also, a tree described in (*) does not contain -1 in the sequence C_1 . For a tree described in (**) the vertex -1 appears in C_1 only in the first place, and the last entry of C_0 (the root of such a tree) is not from V_1 . Therefore, in the code of such a tree there exists $v \in V \setminus V_1$ that appears in C_0 as $U(v'_1)$. We replace this vertex v with -1 and obtain the bijection between generating simplices of $\Delta(\vec{K}_{n_1, n_2, \dots, n_k})$ and directed trees of $\vec{K}_{n_1, n_2, \dots, n_k}$ in which -1 does not occur in C_1 and 1 does not occur at all. For a tree described in (*) the code remains unchanged. The number of these trees is

$$(n-1)^{k-1}(n-n_1-1)^{n_1-1}(n-n_2-1)^{n_2-1} \dots (n-n_k-1)^{n_k-1}. \quad \square$$

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