

Substitutions over infinite alphabet generating $(-\beta)$ -integers

Daniel Dombek

Department of Mathematics FNSPE
Czech Technical University in Prague
Czech Republic
dombedan@fjfi.cvut.cz

1 Introduction

This contribution is devoted to the study of positional numeration systems with negative base introduced by Ito and Sadahiro in 2009, called $(-\beta)$ -expansions. We give an admissibility criterion for more general case of $(-\beta)$ -expansions and discuss the properties of the set of $(-\beta)$ -integers, denoted by $\mathbb{Z}_{-\beta}$. We give a description of distances within $\mathbb{Z}_{-\beta}$ and show that this set can be coded by an infinite word over an infinite alphabet, which is a fixed point of a non-erasing non-trivial morphism.

2 Numeration with negative base

In 1957, Rényi introduced positional numeration system with positive real base $\beta > 1$ (see [7]). The β -expansion of $x \in [0, 1)$ is defined as the digit string $d_\beta(x) = 0 \bullet x_1 x_2 x_3 \dots$, where

$$x_i = \lfloor \beta T_\beta^{i-1}(x) \rfloor \quad \text{and} \quad T_\beta(x) = \beta x - \lfloor \beta x \rfloor.$$

It holds that

$$x = \frac{x_1}{\beta} + \frac{x_2}{\beta^2} + \frac{x_3}{\beta^3} + \dots.$$

Note that this definition can be naturally extended so that any real number has a unique β -expansion, which is usually denoted $d_\beta(x) = x_k x_{k-1} \dots x_1 x_0 \bullet x_{-1} x_{-2} \dots$, where \bullet , the fractional point, separates negative and non-negative powers of β . In analogy with standard integer base, the set \mathbb{Z}_β of β -integers is defined as the set of real numbers having the β -expansion of the form $d_\beta(x) = x_k x_{k-1} \dots x_1 x_0 \bullet 0^\omega$.

$(-\beta)$ -expansions, a numeration system built in analogy with Rényi β -expansions, was introduced in 2009 by Ito and Sadahiro (see [5]). They gave a lexicographic criterion for deciding whether some digit string is the $(-\beta)$ -expansion of some x and also described several properties of $(-\beta)$ -expansions concerning symbolic dynamics and ergodic theory. Note that dynamical properties of $(-\beta)$ -expansions were also studied by Frougny and Lai (see [4]). We take the liberty of defining $(-\beta)$ -expansions in a more general way, while an analogy with positive base numeration can still be easily seen.

Definition 1. Let $-\beta < -1$ be a base and consider $x \in [l, l+1)$, where $l \in \mathbb{R}$ is arbitrary fixed. We define the $(-\beta)$ -expansion of x as the digit string $d(x) = x_1 x_2 x_3 \dots$, with digits x_i given by

$$x_i = \lfloor -\beta T^{i-1}(x) - l \rfloor, \tag{1}$$

where $T(x)$ stands for the generalised $(-\beta)$ -transformation

$$T : [l, l+1) \rightarrow [l, l+1), \quad T(x) = -\beta x - \lfloor -\beta x - l \rfloor. \tag{2}$$

It holds that

$$x = \frac{x_1}{-\beta} + \frac{x_2}{(-\beta)^2} + \frac{x_3}{(-\beta)^3} + \dots$$

and the fractional point is again used in the notation, $d(x) = 0 \bullet x_1 x_2 x_3 \dots$.

The set of digits used in $(-\beta)$ -expansions of numbers (in the latter referred to as the alphabet of $(-\beta)$ -expansions) depends on the choice of l and can be calculated directly from (1) as

$$\mathcal{A}_{-\beta, l} = \{\lfloor -l(\beta + 1) - \beta \rfloor, \dots, \lfloor -l(\beta + 1) \rfloor\}. \quad (3)$$

We may demand that the numeration system possesses various properties. Let us summarise the most natural ones:

- The most common requirement is that zero is an allowed digit. We see that $0 \in \mathcal{A}_{-\beta, l}$ is equivalent to $0 \in [l, l + 1)$ and consequently $l \in (-1, 0]$. Note that this implies $d(0) = 0 \bullet 0^\omega$.
- We may require that $\mathcal{A}_{-\beta, l} = \{0, 1, \dots, \lfloor \beta \rfloor\}$. This is equivalent to the choice $l \in \left(-\frac{\lfloor \beta \rfloor + 1}{\beta + 1}, -\frac{\beta}{\beta + 1}\right]$.
- So far, $(-\beta)$ -expansions were defined only for numbers from $[l, l + 1)$. In Rényi numeration, the β -expansion of arbitrary $x \in \mathbb{R}^+$ (expansions of negative numbers differ only by “-” sign) is defined as $d_\beta(x) = x_k x_{k-1} \dots x_1 x_0 \bullet x_{-1} x_{-2} \dots$, where $k \in \mathbb{N}$ satisfies $\frac{x}{\beta^k} \in [l, l + 1)$ and $d_\beta\left(\frac{x}{\beta^k}\right) = 0 \bullet x_k x_{k-1} x_{k-2} \dots$. The same procedure does not work for $(-\beta)$ -expansions in general. A necessary and sufficient condition for the existence of unique $d(x)$ for all $x \in \mathbb{R}$ is that $-\frac{1}{\beta}[l, l + 1) \subset [l, l + 1)$. This is equivalent to the choice $l \in \left(-\frac{\beta}{\beta + 1}, -\frac{1}{\beta + 1}\right]$. Note that this choice is disjoint with the previous one, so one cannot have uniqueness of $(-\beta)$ -expansions and non-negative digits bounded by β at the same time.

Let us stress that in the following we will need 0 to be a valid digit. Therefore, we shall always assume $l \in (-1, 0]$. Note that we may easily derive that the digits in the alphabet $\mathcal{A}_{-\beta, l}$ are then bounded by $\lceil \beta \rceil$ in modulus.

3 Admissibility

In Rényi numeration there is a natural correspondence between ordering on real numbers and lexicographic ordering on their β -expansions. In $(-\beta)$ -expansions, standard lexicographic ordering is not suitable anymore, hence a different ordering on digit strings is needed.

The so-called alternate order was used in the admissibility condition by Ito and Sadahiro and it will work also in the general case. Let us recall the definition. For the strings

$$u, v \in (\mathcal{A}_{-\beta, l})^{\mathbb{N}}, \quad u = u_1 u_2 u_3 \dots \quad \text{and} \quad v = v_1 v_2 v_3 \dots$$

we say that $u \prec_{alt} v$ (u is less than v in the alternate order) if $u_m(-1)^m < v_m(-1)^m$, where $m = \min\{k \in \mathbb{N} \mid u_k \neq v_k\}$. Note that standard ordering between reals in $[l, l + 1)$ corresponds to the alternate order on their respective $(-\beta)$ -expansions.

Definition 2. An infinite string $x_1 x_2 x_3 \dots$ of integers is called $(-\beta)$ -admissible (or just admissible), if there exists an $x \in [l, l + 1)$ such that $x_1 x_2 x_3 \dots$ is its $(-\beta)$ -expansion, i.e. $x_1 x_2 x_3 \dots = d(x)$.

We give the criterion for $(-\beta)$ -admissibility (proven in [2]) in a form similar to both Parry lexicographic condition (see [6]) and Ito-Sadahiro admissibility criterion (see [5]).

Theorem 3. ([2]) *An infinite string $x_1x_2x_3\cdots$ of integers is $(-\beta)$ -admissible, if and only if*

$$l_1l_2l_3\cdots \preceq_{alt} x_i x_{i+1} x_{i+2} \cdots \prec_{alt} r_1 r_2 r_3 \cdots, \quad \text{for all } i \geq 1, \quad (4)$$

where $l_1l_2l_3\cdots = d(l)$ and $r_1r_2r_3\cdots = d^*(l+1) = \lim_{\varepsilon \rightarrow 0^+} d(l+1-\varepsilon)$.

Remark 4. *Ito and Sadahiro have described the admissibility condition for their numeration system considered with $l = -\frac{\beta}{\beta+1}$. This choice imply for any β the alphabet of the form $\mathcal{A}_{-\beta,l} = \{0, 1, \dots, \lfloor \beta \rfloor\}$. They have shown that in this case the reference strings used in the condition in Theorem 3 (i.e. $d(l) = l_1l_2l_3\cdots$ and $d^*(l+1) = r_1r_2r_3\cdots$) are related in the following way:*

$$r_1r_2r_3\cdots = 0l_1l_2l_3\cdots$$

if $d(l)$ is not purely periodic with odd period length, and,

$$r_1r_2r_3\cdots = (0l_1l_2\cdots l_{q-1}(l_q-1))^\omega,$$

if $d(l) = (l_1l_2\cdots l_q)^\omega$, where q is odd.

Remark 5. *Besides Ito-Sadahiro case and the general one, we may consider another interesting example, the choice $l = -\frac{1}{2}$, $\beta \notin 2\mathbb{Z} + 1$. This leads to a numeration defined on “almost symmetric” interval $[-\frac{1}{2}, \frac{1}{2})$ with symmetric alphabet*

$$\mathcal{A}_{-\beta, -\frac{1}{2}} = \left\{ \left\lfloor \frac{\beta+1}{2} \right\rfloor, \dots, \bar{1}, 0, 1, \dots, \left\lceil \frac{\beta+1}{2} \right\rceil \right\}.$$

Note that we use the notation $(-a) = \bar{a}$ for shorter writing of negative digits. If we denote the reference strings as usual, i.e. $d(-\frac{1}{2}) = l_1l_2l_3\cdots$ and $d^*(\frac{1}{2}) = r_1r_2r_3\cdots$, the following relation can be shown:

$$r_1r_2r_3\cdots = \overline{l_1l_2l_3\cdots}$$

if $d(l)$ is not purely periodic with odd period length, and,

$$r_1r_2r_3\cdots = (\overline{l_1l_2\cdots l_{q-1}(l_q-1)} l_1l_2\cdots l_{q-1}(l_q-1))^\omega,$$

if $d(l) = (l_1l_2\cdots l_q)^\omega$, where q is odd.

4 $(-\beta)$ -integers

We have already discussed basic properties of $(-\beta)$ -expansions and the question of admissibility of digit strings. In the following, $(-\beta)$ -admissibility will be used to define the set of $(-\beta)$ -integers.

Let us define a “value function” γ . Consider a finite digit string $x_{k-1}\cdots x_1x_0$, then $\gamma(x_{k-1}, \dots, x_1x_0) = \sum_{i=0}^{k-1} x_i(-\beta)^i$.

Definition 6. *We call $x \in \mathbb{R}$ a $(-\beta)$ -integer, if there exists a $(-\beta)$ -admissible digit string $x_kx_{k-1}\cdots x_00^\omega$ such that $d(x) = x_kx_{k-1}\cdots x_1x_0 \bullet 0^\omega$. The set of $(-\beta)$ -integers is then defined as*

$$\mathbb{Z}_{-\beta} = \{x \in \mathbb{R} \mid x = \gamma(a_{k-1}a_{k-2}\cdots a_1a_0), a_{k-1}a_{k-2}\cdots a_1a_00^\omega \text{ is } (-\beta)\text{-admissible},$$

or equivalently

$$\mathbb{Z}_{-\beta} = \bigcup_{i \geq 0} (-\beta)^i T^{-i}(0).$$

Note that $(-\beta)$ -expansions of real numbers are not necessarily unique. As was said before, uniqueness holds if and only if $l \in (-\frac{\beta}{\beta+1}, -\frac{1}{\beta+1}]$. Let us demonstrate this ambiguity on the following example.

Example 7. Let β be the greater root of the polynomial $x^2 - 2x - 1$, i.e. $\beta = 1 + \sqrt{2}$, and let $[l, l+1) = [-\frac{\beta^9}{\beta^9+1}, \frac{1}{\beta^9+1})$. Note that $[l, l+1)$ is not invariant under division by $(-\beta)$.

If we want to find the $(-\beta)$ -expansion of number $x \notin [l, l+1)$, we have to find such $k \in \mathbb{N}$ that $\frac{x}{(-\beta)^k} \in [l, l+1)$, compute $d(\frac{x}{(-\beta)^k})$ by definition and then shift the fractional point by k positions to the right. The problem is that, in general, different choices of the exponent k may give different $(-\beta)$ -admissible digit strings which all represent the same number x .

Let us find possible $(-\beta)$ -expansions of 1. It can be shown that $\frac{1}{(-\beta)^k} \in [l, l+1)$ if and only if $k \in \mathbb{N} \setminus \{0, 2, 4, 6, 8\}$ and there are 5 $(-\beta)$ -admissible digit strings representing 1, computed from $(-\beta)$ -expansions of $\frac{1}{(-\beta)^k}$ for $k = 1, 3, 5, 7, 9$ respectively:

$$1 \bullet 0^\omega = 120 \bullet 0^\omega = 13210 \bullet 0^\omega = 1322210 \bullet 0^\omega = 132222210 \bullet 0^\omega.$$

Let us mention some straightforward observations on the properties of $\mathbb{Z}_{-\beta}$:

- $\mathbb{Z}_{-\beta}$ is nonempty if and only if $0 \in \mathcal{A}_{-\beta, l}$, i.e. if and only if $l \in (-1, 0]$.
- The definition implies $-\beta\mathbb{Z}_{-\beta} \subset \mathbb{Z}_{-\beta}$.
- A phenomenon unseen in Rényi numeration arises, there are cases when the set of $(-\beta)$ -integers is trivial, i.e. when $\mathbb{Z}_{-\beta} = \{0\}$. This happens if and only if both numbers $\frac{1}{\beta}$ and $-\frac{1}{\beta}$ are outside of the interval $[l, l+1)$. This can be reformulated as

$$\mathbb{Z}_{-\beta} = \{0\} \Leftrightarrow \beta < -\frac{1}{l} \text{ and } \beta \leq \frac{1}{l+1},$$

and it can be seen that the strictest limitation for β arises when $l = -\frac{1}{2}$. This implies for any choice of $l \in \mathbb{R}$:

$$\mathbb{Z}_{-\beta} \neq \emptyset \text{ and } \beta \geq 2 \Rightarrow \mathbb{Z}_{-\beta} \supsetneq \{0\}.$$

- It holds that $\mathbb{Z}_{-\beta} = \mathbb{Z}$ if and only if $\beta \in \mathbb{N}$.

Remark 8. As was shown in Example 7, in a completely general case of $(-\beta)$ -expansions, there is a problem with ambiguity. Because of this, in the following we shall limit ourselves to the choice $l \in [-\frac{\beta}{\beta+1}, -\frac{1}{\beta+1}]$. Note that we allow Ito-Sadahiro case $l = -\frac{\beta}{\beta+1}$, which also contains ambiguities, but only in countably many cases, which can be avoided by introducing a notion of strong $(-\beta)$ -admissibility.

Definition 9. Let $x_1x_2x_3 \cdots \in \mathcal{A}_{-\beta, l}$. We say that

$$x_1x_2x_3 \cdots \text{ is strongly } (-\beta)\text{-admissible} \quad \text{if} \quad 0x_1x_2x_3 \cdots \text{ is } (-\beta)\text{-admissible}.$$

Remark 10. Note that if $l \in (-\frac{\beta}{\beta+1}, -\frac{1}{\beta+1}]$, the notions of strong admissibility and admissibility coincide. In the case $l = -\frac{\beta}{\beta+1}$, the only numbers with non-unique expansions are those of the form $(-\beta)^k l$, which have exactly two possible expansions using digit strings $l_1l_2l_3 \cdots$ and $1l_1l_2l_3 \cdots$. While both are $(-\beta)$ -admissible, only the latter is also strongly $(-\beta)$ -admissible.

In order to describe distances between adjacent $(-\beta)$ -integers, we will study ordering of finite digit strings in the alternate order. Denote by $\mathcal{S}(k)$ the set of infinite $(-\beta)$ -admissible digit strings such that erasing a prefix of length k yields 0^ω , i.e. for $k \geq 0$, we have

$$\mathcal{S}(k) = \{a_{k-1}a_{k-2}\cdots a_00^\omega \mid a_{k-1}a_{k-2}\cdots a_00^\omega \text{ is } (-\beta)\text{-admissible}\},$$

in particular $\mathcal{S}(0) = \{0^\omega\}$. For a fixed k , the set $\mathcal{S}(k)$ is finite. Denote by $\text{Max}(k)$ the string $a_{k-1}a_{k-2}\cdots a_00^\omega$ which is maximal in $\mathcal{S}(k)$ with respect to the alternate order and by $\text{max}(k)$ its prefix of length k , i.e. $\text{Max}(k) = \text{max}(k)0^\omega$. Similarly, we define $\text{Min}(k)$ and $\text{min}(k)$. Thus,

$$\text{Min}(k) \preceq_{alt} r \preceq_{alt} \text{Max}(k), \quad \text{for all digit strings } r \in \mathcal{S}(k).$$

With this notation we can give a theorem describing distances in $\mathbb{Z}_{-\beta}$ valid for cases $l \in [-\frac{\beta}{\beta+1}, -\frac{1}{\beta+1}]$. Note that for case $l = -\frac{\beta}{\beta+1}$ it was proven in [1].

Theorem 11. *Let $x < y$ be two consecutive $(-\beta)$ -integers. Then there exist a finite string w over the alphabet $\mathcal{A}_{-\beta,l}$, a non-negative integer $k \in \{0, 1, 2, \dots\}$ and a positive digit $d \in \mathcal{A}_{-\beta,l} \setminus \{0\}$ such that $w(d-1)\text{Max}(k)$ and $w\text{dMin}(k)$ are strongly $(-\beta)$ -admissible strings and*

$$\begin{aligned} x = \gamma(w(d-1)\text{max}(k)) &< y = \gamma(w\text{dmin}(k)) && \text{for } k \text{ even,} \\ x = \gamma(w\text{dmin}(k)) &< y = \gamma(w(d-1)\text{max}(k)) && \text{for } k \text{ odd.} \end{aligned}$$

In particular, the distance $y - x$ between these $(-\beta)$ -integers depends only on k and equals to

$$\Delta_k := \left| (-\beta)^k + \gamma(\text{min}(k)) - \gamma(\text{max}(k)) \right|. \quad (5)$$

5 Coding $\mathbb{Z}_{-\beta}$ by an infinite word

Note that in order to get an explicit formula for distances from Theorem 3, knowledge of reference strings $\text{min}(k)$ and $\text{max}(k)$ is necessary. These depend on both reference strings $d(l)$ and $d^*(l+1)$. Concerning the form of $\text{min}(k)$ and $\text{max}(k)$ we provide the following proposition.

Proposition 12. *Let $\beta > 1$. Denote $d(l) = l_1l_2l_3\cdots$, $d^*(l+1) = r_1r_2r_3\cdots$.*

- $\text{min}(0) = \text{max}(0) = \varepsilon$,
- for $k \geq 1$ either $\text{min}(k) = l_1l_2\cdots l_k$ or there exists $m(k) \in \{0, \dots, k-1\}$ such that

$$\text{min}(k) = \begin{cases} l_1l_2\cdots(l_{k-m(k)}+1)\text{min}(m(k)) & \text{if } k-m(k) \text{ even} \\ l_1l_2\cdots(l_{k-m(k)}-1)\text{max}(m(k)) & \text{if } k-m(k) \text{ odd} \end{cases}$$

- for $k \geq 1$ either $\text{max}(k) = r_1r_2\cdots r_k$ or there exists $m'(k) \in \{0, \dots, k-1\}$ such that

$$\text{max}(k) = \begin{cases} r_1r_2\cdots(r_{k-m'(k)}-1)\text{max}(m'(k)) & \text{if } k-m'(k) \text{ even} \\ r_1r_2\cdots(r_{k-m'(k)}+1)\text{min}(m'(k)) & \text{if } k-m'(k) \text{ odd} \end{cases}$$

Computing $\min(k)$ and $\max(k)$ for a general choice of l may lead to difficult discussion, however, in special cases an important relation between $d(l)$ and $d^*(l+1)$ arises and eases the computation. Examples were given in Remarks 4 and 5.

Let us now describe how we can code the set of $(-\beta)$ -integers by an infinite word over the infinite alphabet \mathbb{N} .

Let $(z_n)_{n \in \mathbb{Z}}$ be a strictly increasing sequence satisfying

$$z_0 = 0 \quad \text{and} \quad \mathbb{Z}_{-\beta} = \{z_n \mid n \in \mathbb{Z}\}.$$

We define a bidirectional infinite word over an infinite alphabet $\mathbf{v}_{-\beta} \in \mathbb{N}^{\mathbb{Z}}$, which codes the set of $(-\beta)$ -integers. According to Theorem 11, for any $n \in \mathbb{Z}$ there exist a unique $k \in \mathbb{N}$, a word w with prefix 0 and a letter d such that

$$z_{n+1} - z_n = |\gamma(w(d-1)\max(k)) - \gamma(wd\min(k))|.$$

We define the word $\mathbf{v}_{-\beta} = (v_i)_{i \in \mathbb{Z}}$ by $v_n = k$.

Theorem 13. *Let $\mathbf{v}_{-\beta}$ be the word associated with $(-\beta)$ -integers. There exists an antimorphism $\Phi : \mathbb{N}^* \rightarrow \mathbb{N}^*$ such that $\Psi = \Phi^2$ is a non-erasing non-identical morphism and $\Psi(\mathbf{v}_{-\beta}) = \mathbf{v}_{-\beta}$. Φ is always of the form*

$$\Phi(2l) = S_{2l}(2l+1)\widetilde{R}_{2l} \quad \text{and} \quad \Phi(2l+1) = R_{2l+1}(2l+2)\widetilde{S}_{2l+1},$$

where \widetilde{u} denotes the reversal of the word u and words R_j, S_j depend only on j and on $\min(k), \max(k)$ with $k \in \{j, j+1\}$.

The proof is based on the self-similarity of $\mathbb{Z}_{-\beta}$, i.e. $-\beta\mathbb{Z}_{-\beta} \subset \mathbb{Z}_{-\beta}$, and on the following idea. Let $x = \gamma(w(d-1)\max(k)) < y = \gamma(wd\min(k))$ be two neighbours in $\mathbb{Z}_{-\beta}$ with gap Δ_k and suppose only k even. If we multiply both x and y by $(-\beta)$, we get a longer gap with possibly more $(-\beta)$ -integers in between. It can be shown that between $-\beta y$ and $-\beta x$ there is always a gap Δ_{k+1} . Hence the description is of the form $\Phi(k) = S_k(k+1)\widetilde{R}_k$, where the word S_k codes the distances between $(-\beta)$ -integers in $[\gamma(wd\min(k)0), \gamma(wd\min(k+1))]$ and, similarly, R_k encodes distances within the interval $[\gamma(w(d-1)\max(k)0), \gamma(w(d-1)\max(k+1))]$.

As it turns out, in some cases (mostly when reference strings $l_1 l_2 l_3 \dots$ and $r_1 r_2 r_3 \dots$ are eventually periodic of a particular form) we can find a letter-to-letter projection to a finite alphabet $\Pi : \mathbb{N} \rightarrow \mathcal{B}$ with $\mathcal{B} \subset \mathbb{N}$, such that $\mathbf{u}_{-\beta} = \Pi \mathbf{v}_{-\beta}$ also encodes $\mathbb{Z}_{-\beta}$ and it is a fixed point of a an antimorphism $\varphi = \Pi \circ \Phi$ over the finite alphabet \mathcal{B} . Clearly, the square of φ is then a non-erasing morphism over \mathcal{B} which fixes $\mathbf{u}_{-\beta}$.

Let us mention that $(-\beta)$ -integers in the Ito-Sadahiro case $l = -\frac{\beta}{\beta+1}$ are also subject of [8]. For β with eventually periodic $d(l)$, Steiner finds a coding of $\mathbb{Z}_{-\beta}$ by a finite alphabet and shows, using only the properties of the $(-\beta)$ -transformation, that the word is a fixed point of a non-trivial morphism. Our approach is of a combinatorial nature, follows a similar idea as in [1] and shows existence of an antimorphism for any base β .

To illustrate the results, let us conclude this contribution by an example.

Example 14. *Let β be the real root of $x^3 - 3x^2 - 4x - 2$ (β Pisot, ≈ 4.3) and $l = -\frac{1}{2}$. The admissibility condition gives us for any admissible digit string $(x_i)_{i \geq 0}$:*

$$201^\omega \preceq_{alt} x_i x_{i+1} x_{i+2} \dots \prec_{alt} \overline{201}^\omega \quad \text{for all } x \geq 0.$$

We obtain

$$\min(0) = \varepsilon, \quad \min(1) = 2, \quad \min(2) = 20$$

and

$$\min(2k+1) = 20(11)^{k-1}0, \quad \min(2k+2) = 20(11)^k \quad \text{for } k \geq 1.$$

Clearly it holds that $\max(i) = \overline{\min(i)}$ for all $i \in \mathbb{N}$.

Theorem 11 gives us the following distances within $\mathbb{Z}_{-\beta}$:

$$\Delta_0 = 1, \quad \Delta_1 = -1 + \frac{4}{\beta} + \frac{2}{\beta^2}, \quad \text{and} \quad \Delta_{2k} = 1 - \frac{2}{\beta} - \frac{2}{\beta^2}, \quad \Delta_{2k+1} = 1 + \frac{2}{\beta} + \frac{2}{\beta^2} \quad \text{for } k \geq 1.$$

Finally, the antimorphism $\Phi : \mathbb{N}^* \rightarrow \mathbb{N}^*$ is given by

$$\begin{aligned} 0 &\rightarrow 0^2 10^2, \\ 1 &\rightarrow 2, \\ 2 &\rightarrow 3, \end{aligned}$$

and for $k \geq 1$

$$\begin{aligned} 2k+1 &\rightarrow 0^2 10(2k+2)010^2, \\ 2k+2 &\rightarrow 2k+3. \end{aligned}$$

It can be easily seen that a projection from \mathbb{N} to a finite alphabet exists and a final antimorphism $\varphi : \{0, 1, 2, 3\}^* \rightarrow \{0, 1, 2, 3\}^*$ is of the form

$$\begin{aligned} 0 &\rightarrow 0^2 10^2, \\ 1 &\rightarrow 2, \\ 2 &\rightarrow 3, \\ 3 &\rightarrow 0^2 102010^2. \end{aligned}$$

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