

GNU Generic Security Service Library

GSS-API Library for the GNU system
for version 0.0.19, 9 January 2007

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1 Introduction

GSS is an implementation of the Generic Security Service Application Program Interface (GSS-API). GSS-API is used by network servers to provide security services, e.g., to authenticate SMTP/IMAP clients against SMTP/IMAP servers. GSS consists of a library and a manual.

GSS is developed for the GNU/Linux system, but runs on over 20 platforms including most major Unix platforms and Windows, and many kind of devices including iPAQ handhelds and S/390 mainframes.

GSS is a GNU project, and is licensed under the GNU General Public License.

1.1 Getting Started

This manual documents the GSS programming interface. All functions and data types provided by the library are explained.

The reader is assumed to possess basic familiarity with GSS-API and network programming in C or C++. For general GSS-API information, and some programming examples, there is a guide available online at <http://docs.sun.com/db/doc/816-1331>.

This manual can be used in several ways. If read from the beginning to the end, it gives a good introduction into the library and how it can be used in an application. Forward references are included where necessary. Later on, the manual can be used as a reference manual to get just the information needed about any particular interface of the library. Experienced programmers might want to start looking at the examples at the end of the manual, and then only read up those parts of the interface which are unclear.

1.2 Features

GSS might have a couple of advantages over other libraries doing a similar job.

It's Free Software

Anybody can use, modify, and redistribute it under the terms of the GNU General Public License.

It's thread-safe

No global variables are used and multiple library handles and session handles may be used in parallel.

It's internationalized

It handles non-ASCII names and user visible strings used in the library (e.g., error messages) can be translated into the users' language.

It's portable

It should work on all Unix like operating systems, including Windows.

1.3 GSS-API Overview

This section describes GSS-API from a protocol point of view.

The Generic Security Service Application Programming Interface provides security services to calling applications. It allows a communicating application to authenticate the user

associated with another application, to delegate rights to another application, and to apply security services such as confidentiality and integrity on a per-message basis.

There are four stages to using the GSS-API:

1. The application acquires a set of credentials with which it may prove its identity to other processes. The application's credentials vouch for its global identity, which may or may not be related to any local username under which it may be running.
2. A pair of communicating applications establish a joint security context using their credentials. The security context is a pair of GSS-API data structures that contain shared state information, which is required in order that per-message security services may be provided. Examples of state that might be shared between applications as part of a security context are cryptographic keys, and message sequence numbers. As part of the establishment of a security context, the context initiator is authenticated to the responder, and may require that the responder is authenticated in turn. The initiator may optionally give the responder the right to initiate further security contexts, acting as an agent or delegate of the initiator. This transfer of rights is termed delegation, and is achieved by creating a set of credentials, similar to those used by the initiating application, but which may be used by the responder.

To establish and maintain the shared information that makes up the security context, certain GSS-API calls will return a token data structure, which is an opaque data type that may contain cryptographically protected data. The caller of such a GSS-API routine is responsible for transferring the token to the peer application, encapsulated if necessary in an application- application protocol. On receipt of such a token, the peer application should pass it to a corresponding GSS-API routine which will decode the token and extract the information, updating the security context state information accordingly.

3. Per-message services are invoked to apply either: integrity and data origin authentication, or confidentiality, integrity and data origin authentication to application data, which are treated by GSS-API as arbitrary octet-strings. An application transmitting a message that it wishes to protect will call the appropriate GSS-API routine (`gss_get_mic` or `gss_wrap`) to apply protection, specifying the appropriate security context, and send the resulting token to the receiving application. The receiver will pass the received token (and, in the case of data protected by `gss_get_mic`, the accompanying message-data) to the corresponding decoding routine (`gss_verify_mic` or `gss_unwrap`) to remove the protection and validate the data.
4. At the completion of a communications session (which may extend across several transport connections), each application calls a GSS-API routine to delete the security context. Multiple contexts may also be used (either successively or simultaneously) within a single communications association, at the option of the applications.

1.4 Supported Platforms

GSS has at some point in time been tested on the following platforms.

1. Debian GNU/Linux 3.0 (Woody)

GCC 2.95.4 and GNU Make. This is the main development platform. `alphaev67-unknown-linux-gnu`, `alphaev6-unknown-linux-gnu`, `arm-unknown-linux-gnu`, `hppa-unknown-linux-gnu`, `hppa64-unknown-linux-gnu`, `i686-pc-linux-gnu`,

- `ia64-unknown-linux-gnu`, `m68k-unknown-linux-gnu`, `mips-unknown-linux-gnu`,
`mipsel-unknown-linux-gnu`, `powerpc-unknown-linux-gnu`, `s390-ibm-linux-gnu`,
`sparc-unknown-linux-gnu`.
2. Debian GNU/Linux 2.1
GCC 2.95.1 and GNU Make. `armv4l-unknown-linux-gnu`.
 3. Tru64 UNIX
Tru64 UNIX C compiler and Tru64 Make. `alphaev67-dec-osf5.1`, `alphaev68-dec-osf5.1`.
 4. SuSE Linux 7.1
GCC 2.96 and GNU Make. `alphaev6-unknown-linux-gnu`, `alphaev67-unknown-linux-gnu`.
 5. SuSE Linux 7.2a
GCC 3.0 and GNU Make. `ia64-unknown-linux-gnu`.
 6. RedHat Linux 7.2
GCC 2.96 and GNU Make. `alphaev6-unknown-linux-gnu`, `alphaev67-unknown-linux-gnu`, `ia64-unknown-linux-gnu`.
 7. RedHat Linux 8.0
GCC 3.2 and GNU Make. `i686-pc-linux-gnu`.
 8. RedHat Advanced Server 2.1
GCC 2.96 and GNU Make. `i686-pc-linux-gnu`.
 9. Slackware Linux 8.0.01
GCC 2.95.3 and GNU Make. `i686-pc-linux-gnu`.
 10. Mandrake Linux 9.0
GCC 3.2 and GNU Make. `i686-pc-linux-gnu`.
 11. IRIX 6.5
MIPS C compiler, IRIX Make. `mips-sgi-irix6.5`.
 12. AIX 4.3.2
IBM C for AIX compiler, AIX Make. `rs6000-ibm-aix4.3.2.0`.
 13. Microsoft Windows 2000 (Cygwin)
GCC 3.2, GNU make. `i686-pc-cygwin`.
 14. HP-UX 11
HP-UX C compiler and HP Make. `ia64-hp-hpux11.22`, `hppa2.0w-hp-hpux11.11`.
 15. SUN Solaris 2.8
Sun WorkShop Compiler C 6.0 and SUN Make. `sparc-sun-solaris2.8`.
 16. NetBSD 1.6
GCC 2.95.3 and GNU Make. `alpha-unknown-netbsd1.6`, `i386-unknown-netbsdelf1.6`.
 17. OpenBSD 3.1 and 3.2
GCC 2.95.3 and GNU Make. `alpha-unknown-openbsd3.1`, `i386-unknown-openbsd3.1`.

18. FreeBSD 4.7
GCC 2.95.4 and GNU Make. `alpha-unknown-freebsd4.7`, `i386-unknown-freebsd4.7`.
19. Cross compiled to uClinux/uClibc on Motorola Coldfire.
GCC 3.4 and GNU Make `m68k-uclinux-elf`.

If you use GSS on, or port GSS to, a new platform please report it to the author.

1.5 Commercial Support

Commercial support is available for users of GNU GSS. The kind of support that can be purchased may include:

- Implement new features. Such as a new GSS-API mechanism.
- Port GSS to new platforms. This could include porting to an embedded platforms that may need memory or size optimization.
- Integrating GSS as a security environment in your existing project.
- System design of components related to GSS-API.

If you are interested, please write to:

Simon Josefsson Datakonsult
Hagagatan 24
113 47 Stockholm
Sweden

E-mail: simon@josefsson.org

If your company provide support related to GNU GSS and would like to be mentioned here, contact the author (see [Section 1.7 \[Bug Reports\]](#), page 5).

1.6 Downloading and Installing

The package can be downloaded from several places, including:

<http://josefsson.org/gss/releases/>

The latest version is stored in a file, e.g., `gss-0.0.19.tar.gz` where the `0.0.19` indicate the highest version number.

The package is then extracted, configured and built like many other packages that use Autoconf. For detailed information on configuring and building it, refer to the `INSTALL` file that is part of the distribution archive.

Here is an example terminal session that download, configure, build and install the package. You will need a few basic tools, such as `sh`, `make` and `cc`.

```
$ wget -q http://josefsson.org/gss/releases/gss-0.0.19.tar.gz
$ tar xzf gss-0.0.19.tar.gz
$ cd gss-0.0.19/
$ ./configure
...
$ make
...
```



```
$ make install
...
```

After that GSS should be properly installed and ready for use.

1.7 Bug Reports

If you think you have found a bug in GSS, please investigate it and report it.

- Please make sure that the bug is really in GSS, and preferably also check that it hasn't already been fixed in the latest version.
- You have to send us a test case that makes it possible for us to reproduce the bug.
- You also have to explain what is wrong; if you get a crash, or if the results printed are not good and in that case, in what way. Make sure that the bug report includes all information you would need to fix this kind of bug for someone else.

Please make an effort to produce a self-contained report, with something definite that can be tested or debugged. Vague queries or piecemeal messages are difficult to act on and don't help the development effort.

If your bug report is good, we will do our best to help you to get a corrected version of the software; if the bug report is poor, we won't do anything about it (apart from asking you to send better bug reports).

If you think something in this manual is unclear, or downright incorrect, or if the language needs to be improved, please also send a note.

Send your bug report to:

`'bug-gss@josefsson.org'`

1.8 Contributing

If you want to submit a patch for inclusion – from solve a typo you discovered, up to adding support for a new feature – you should submit it as a bug report (see [Section 1.7 \[Bug Reports\], page 5](#)). There are some things that you can do to increase the chances for it to be included in the official package.

Unless your patch is very small (say, under 10 lines) we require that you assign the copyright of your work to the Free Software Foundation. This is to protect the freedom of the project. If you have not already signed papers, we will send you the necessary information when you submit your contribution.

For contributions that doesn't consist of actual programming code, the only guidelines are common sense. Use it.

For code contributions, a number of style guides will help you:

- Coding Style. Follow the GNU Standards document (see [\[top\]](#), page [\[undefined\]](#)).

If you normally code using another coding standard, there is no problem, but you should use `'indent'` to reformat the code (see [\[top\]](#), page [\[undefined\]](#)) before submitting your work.

- Use the unified diff format `'diff -u'`.

- Return errors. No reason whatsoever should abort the execution of the library. Even memory allocation errors, e.g. when malloc return NULL, should work although result in an error code.
- Design with thread safety in mind. Don't use global variables. Don't even write to per-handle global variables unless the documented behaviour of the function you write is to write to the per-handle global variable.
- Avoid using the C math library. It causes problems for embedded implementations, and in most situations it is very easy to avoid using it.
- Document your functions. Use comments before each function headers, that, if properly formatted, are extracted into Texinfo manuals and GTK-DOC web pages.
- Supply a ChangeLog and NEWS entries, where appropriate.

1.9 Planned Features

This is also known as the “todo list”. If you like to start working on anything, please let me know so work duplication can be avoided.

- Support non-blocking mode. This would be an API extension. It could work by forking a process and interface to it, or by using a user-specific daemon. E.g., `h = START(accept_sec_context(...))`, `FINISHED(h)`, `ret = FINISH(h)`, `ABORT(h)`.
- Compute MD5 of non-empty channel bindings.
- Support loadable modules via dlopen, a'la Solaris GSS.
- Port to Cyclone? CCured?

2 Preparation

To use GSS, you have to perform some changes to your sources and the build system. The necessary changes are small and explained in the following sections. At the end of this chapter, it is described how the library is initialized, and how the requirements of the library are verified.

A faster way to find out how to adapt your application for use with GSS may be to look at the examples at the end of this manual.

2.1 Header

All standard interfaces (data types and functions) of the official GSS API are defined in the header file `'gss/api.h'`. The file is taken verbatim from the RFC (after correcting a few typos) where it is known as `'gssapi.h'`. However, to be able to co-exist gracefully with other GSS-API implementation, the name `'gssapi.h'` was changed.

The header file `'gss.h'` includes `'gss/api.h'`, and declares a few non-standard extensions (by including `'gss/ext.h'`), takes care of including header files related to all supported mechanisms (e.g., `'gss/krb5.h'`) and finally adds C++ namespace protection of all definitions. Therefore, including `'gss.h'` in your project is recommended over `'gss/api.h'`. If using `'gss.h'` instead of `'gss/api.h'` causes problems, it should be regarded a bug.

You must include either file in all programs using the library, either directly or through some other header file, like this:

```
#include <gss.h>
```

The name space of GSS is `gss_*` for function names, `gss_*` for data types and `GSS_*` for other symbols. In addition the same name prefixes with one prepended underscore are reserved for internal use and should never be used by an application.

Each supported GSS mechanism may want to expose mechanism specific functionality, and can do so through one or more header files under the `'gss/'` directory. The Kerberos 5 mechanism uses the file `'gss/krb5.h'`, but again, it is included (with C++ namespace fixes) from `'gss.h'`.

2.2 Initialization

GSS does not need to be initialized before it can be used.

In order to take advantage of the internationalisation features in GSS, e.g. translated error messages, the application must set the current locale using `setlocale()` before calling, e.g., `gss_display_status()`. This is typically done in `main()` as in the following example.

```
#include <gss.h>
#include <locale.h>
...
setlocale (LC_ALL, "");
```

2.3 Version Check

It is often desirable to check that the version of GSS used is indeed one which fits all requirements. Even with binary compatibility new features may have been introduced but

due to problem with the dynamic linker an old version is actually used. So you may want to check that the version is okay right after program startup. The function is called `gss_check_version()` and is described formally in See [Chapter 4 \[Extended GSS API\]](#), page 54.

The normal way to use the function is to put something similar to the following early in your `main()`:

```
#include <gss.h>
...
if (!gss_check_version (GSS_VERSION))
{
    printf ("gss_check_version() failed:\n"
           "Header file incompatible with shared library.\n");
    exit(1);
}
```

2.4 Building the source

If you want to compile a source file that includes the `'gss.h'` header file, you must make sure that the compiler can find it in the directory hierarchy. This is accomplished by adding the path to the directory in which the header file is located to the compilers include file search path (via the `'-I'` option).

However, the path to the include file is determined at the time the source is configured. To solve this problem, GSS uses the external package `pkg-config` that knows the path to the include file and other configuration options. The options that need to be added to the compiler invocation at compile time are output by the `'--cflags'` option to `pkg-config gss`. The following example shows how it can be used at the command line:

```
gcc -c foo.c 'pkg-config gss --cflags'
```

Adding the output of `'pkg-config gss --cflags'` to the compilers command line will ensure that the compiler can find the `'gss.h'` header file.

A similar problem occurs when linking the program with the library. Again, the compiler has to find the library files. For this to work, the path to the library files has to be added to the library search path (via the `'-L'` option). For this, the option `'--libs'` to `pkg-config gss` can be used. For convenience, this option also outputs all other options that are required to link the program with the GSS library (for instance, the `'-lshishi'` option). The example shows how to link `'foo.o'` with GSS into a program `foo`.

```
gcc -o foo foo.o 'pkg-config gss --libs'
```

Of course you can also combine both examples to a single command by specifying both options to `pkg-config`:

```
gcc -o foo foo.c 'pkg-config gss --cflags --libs'
```

2.5 Out of Memory handling

The GSS API does not have a standard error code for the out of memory error condition. Instead of adding a non-standard error code, this library has chosen to adopt a different strategy. The rationale for the strategy chosen is that out of memory handling happens in rare situations, but performing the out of memory error handling after many API function

invocations make source code harder to read. That may make it harder to spot more serious problems. The strategy chosen improve code readability and robustness.

For most applications, aborting the application with an error message when the out of memory situation occur is the best that can be wished for. This is how the library behaves by default.

However, we realize that some applications may not want to have the GSS library abort execution in any situation, or that the error message should go somewhere else than to `stderr`.

To meet this need, the GSS library support a callback function to let the application regain control, and perform its own cleanups, when an out of memory situation has occurred. The application can define a function and set the library variable `gss_alloc_fail_function` to that function. This could be implemented as follows.

```
#include <gss/ext.h>
...
void my_gss_alloc_failer (void)
{
    syslog (LOG_CRIT, "GSS library out of memory");
    exit (1);
}
...
int
main (int argc, char *argv[])
{
    ...
    gss_alloc_fail_function = my_gss_alloc_failer;
    ...
}
```

The GSS library will invoke `gss_alloc_fail_function` if an out of memory error occurs. Note that after this all previously allocated GSS library variables are in an undefined state, so you must not use any previously allocated GSS variables again. You could discard all previous GSS variable and start from scratch, though. The hook is only intended to allow the application to log the situation in a special way. Of course, care must be taken to not allocate more memory, as that will likely also fail.

As can be seen, the callback is a global variable, and is thus not thread safe. It is assumed that if `malloc` fail in one thread, it likely do the same for all threads within the application.

3 Standard GSS API

3.1 Simple Data Types

The following conventions are used by the GSS-API C-language bindings:

3.1.1 Integer types

GSS-API uses the following integer data type:

```
OM_uint32    32-bit unsigned integer
```

3.1.2 String and similar data

Many of the GSS-API routines take arguments and return values that describe contiguous octet-strings. All such data is passed between the GSS-API and the caller using the `gss_buffer_t` data type. This data type is a pointer to a buffer descriptor, which consists of a length field that contains the total number of bytes in the datum, and a value field which contains a pointer to the actual datum:

```
typedef struct gss_buffer_desc_struct {
    size_t    length;
    void      *value;
} gss_buffer_desc, *gss_buffer_t;
```

Storage for data returned to the application by a GSS-API routine using the `gss_buffer_t` conventions is allocated by the GSS-API routine. The application may free this storage by invoking the `gss_release_buffer` routine. Allocation of the `gss_buffer_desc` object is always the responsibility of the application; unused `gss_buffer_desc` objects may be initialized to the value `GSS_C_EMPTY_BUFFER`.

3.1.2.1 Opaque data types

Certain multiple-word data items are considered opaque data types at the GSS-API, because their internal structure has no significance either to the GSS-API or to the caller. Examples of such opaque data types are the `input_token` parameter to `gss_init_sec_context` (which is opaque to the caller), and the `input_message` parameter to `gss_wrap` (which is opaque to the GSS-API). Opaque data is passed between the GSS-API and the application using the `gss_buffer_t` datatype.

3.1.2.2 Character strings

Certain multiple-word data items may be regarded as simple ISO Latin-1 character strings. Examples are the printable strings passed to `gss_import_name` via the `input_name_buffer` parameter. Some GSS-API routines also return character strings. All such character strings are passed between the application and the GSS-API implementation using the `gss_buffer_t` datatype, which is a pointer to a `gss_buffer_desc` object.

When a `gss_buffer_desc` object describes a printable string, the length field of the `gss_buffer_desc` should only count printable characters within the string. In particular, a trailing NUL character should NOT be included in the length count, nor should either the GSS-API implementation or the application assume the presence of an uncounted trailing NUL.

3.1.3 Object Identifiers

Certain GSS-API procedures take parameters of the type `gss_OID`, or Object identifier. This is a type containing ISO-defined tree-structured values, and is used by the GSS-API caller to select an underlying security mechanism and to specify namespaces. A value of type `gss_OID` has the following structure:

```
typedef struct gss_OID_desc_struct {
    OM_uint32    length;
    void         *elements;
} gss_OID_desc, *gss_OID;
```

The `elements` field of this structure points to the first byte of an octet string containing the ASN.1 BER encoding of the value portion of the normal BER TLV encoding of the `gss_OID`. The `length` field contains the number of bytes in this value. For example, the `gss_OID` value corresponding to `iso(1) identified-organization(3) icd-ecma(12) member-company(2) dec(1011) cryptoAlgorithms(7) DASS(5)`, meaning the DASS X.509 authentication mechanism, has a `length` field of 7 and an `elements` field pointing to seven octets containing the following octal values: 53,14,2,207,163,7,5. GSS-API implementations should provide constant `gss_OID` values to allow applications to request any supported mechanism, although applications are encouraged on portability grounds to accept the default mechanism. `gss_OID` values should also be provided to allow applications to specify particular name types (see section 3.10). Applications should treat `gss_OID_desc` values returned by GSS-API routines as read-only. In particular, the application should not attempt to deallocate them with `free()`.

3.1.4 Object Identifier Sets

Certain GSS-API procedures take parameters of the type `gss_OID_set`. This type represents one or more object identifiers (see [\[Object Identifiers\]](#), page 11). A `gss_OID_set` object has the following structure:

```
typedef struct gss_OID_set_desc_struct {
    size_t      count;
    gss_OID     elements;
} gss_OID_set_desc, *gss_OID_set;
```

The `count` field contains the number of OIDs within the set. The `elements` field is a pointer to an array of `gss_OID_desc` objects, each of which describes a single OID. `gss_OID_set` values are used to name the available mechanisms supported by the GSS-API, to request the use of specific mechanisms, and to indicate which mechanisms a given credential supports.

All OID sets returned to the application by GSS-API are dynamic objects (the `gss_OID_set_desc`, the "elements" array of the set, and the "elements" array of each member OID are all dynamically allocated), and this storage must be deallocated by the application using the `gss_release_oid_set` routine.

3.2 Complex Data Types

3.2.1 Credentials

A credential handle is a caller-opaque atomic datum that identifies a GSS-API credential data structure. It is represented by the caller-opaque type `gss_cred_id_t`.

GSS-API credentials can contain mechanism-specific principal authentication data for multiple mechanisms. A GSS-API credential is composed of a set of credential-elements, each of which is applicable to a single mechanism. A credential may contain at most one credential-element for each supported mechanism. A credential-element identifies the data needed by a single mechanism to authenticate a single principal, and conceptually contains two credential-references that describe the actual mechanism-specific authentication data, one to be used by GSS-API for initiating contexts, and one to be used for accepting contexts. For mechanisms that do not distinguish between acceptor and initiator credentials, both references would point to the same underlying mechanism-specific authentication data.

Credentials describe a set of mechanism-specific principals, and give their holder the ability to act as any of those principals. All principal identities asserted by a single GSS-API credential should belong to the same entity, although enforcement of this property is an implementation-specific matter. The GSS-API does not make the actual credentials available to applications; instead a credential handle is used to identify a particular credential, held internally by GSS-API. The combination of GSS-API credential handle and mechanism identifies the principal whose identity will be asserted by the credential when used with that mechanism.

The `gss_init_sec_context` and `gss_accept_sec_context` routines allow the value `GSS_C_NO_CREDENTIAL` to be specified as their credential handle parameter. This special credential-handle indicates a desire by the application to act as a default principal.

3.2.2 Contexts

The `gss_ctx_id_t` data type contains a caller-opaque atomic value that identifies one end of a GSS-API security context.

The security context holds state information about each end of a peer communication, including cryptographic state information.

3.2.3 Authentication tokens

A token is a caller-opaque type that GSS-API uses to maintain synchronization between the context data structures at each end of a GSS-API security context. The token is a cryptographically protected octet-string, generated by the underlying mechanism at one end of a GSS-API security context for use by the peer mechanism at the other end. Encapsulation (if required) and transfer of the token are the responsibility of the peer applications. A token is passed between the GSS-API and the application using the `gss_buffer_t` conventions.

3.2.4 Interprocess tokens

Certain GSS-API routines are intended to transfer data between processes in multi-process programs. These routines use a caller-opaque octet-string, generated by the GSS-API in one process for use by the GSS-API in another process. The calling application is responsible for transferring such tokens between processes in an OS-specific manner. Note that, while GSS-API implementors are encouraged to avoid placing sensitive information within interprocess tokens, or to cryptographically protect them, many implementations will be unable to

avoid placing key material or other sensitive data within them. It is the application's responsibility to ensure that interprocess tokens are protected in transit, and transferred only to processes that are trustworthy. An interprocess token is passed between the GSS-API and the application using the `gss_buffer_t` conventions.

3.2.5 Names

A name is used to identify a person or entity. GSS-API authenticates the relationship between a name and the entity claiming the name.

Since different authentication mechanisms may employ different namespaces for identifying their principals, GSSAPI's naming support is necessarily complex in multi-mechanism environments (or even in some single-mechanism environments where the underlying mechanism supports multiple namespaces).

Two distinct representations are defined for names:

- An internal form. This is the GSS-API "native" format for names, represented by the implementation-specific `gss_name_t` type. It is opaque to GSS-API callers. A single `gss_name_t` object may contain multiple names from different namespaces, but all names should refer to the same entity. An example of such an internal name would be the name returned from a call to the `gss_inquire_cred` routine, when applied to a credential containing credential elements for multiple authentication mechanisms employing different namespaces. This `gss_name_t` object will contain a distinct name for the entity for each authentication mechanism.

For GSS-API implementations supporting multiple namespaces, objects of type `gss_name_t` must contain sufficient information to determine the namespace to which each primitive name belongs.

- Mechanism-specific contiguous octet-string forms. A format capable of containing a single name (from a single namespace). Contiguous string names are always accompanied by an object identifier specifying the namespace to which the name belongs, and their format is dependent on the authentication mechanism that employs the name. Many, but not all, contiguous string names will be printable, and may therefore be used by GSS-API applications for communication with their users.

Routines (`gss_import_name` and `gss_display_name`) are provided to convert names between contiguous string representations and the internal `gss_name_t` type. `gss_import_name` may support multiple syntaxes for each supported namespace, allowing users the freedom to choose a preferred name representation. `gss_display_name` should use an implementation-chosen printable syntax for each supported name-type.

If an application calls `gss_display_name`, passing the internal name resulting from a call to `gss_import_name`, there is no guarantee the the resulting contiguous string name will be the same as the original imported string name. Nor do name-space identifiers necessarily survive unchanged after a journey through the internal name-form. An example of this might be a mechanism that authenticates X.500 names, but provides an algorithmic mapping of Internet DNS names into X.500. That mechanism's implementation of `gss_import_name` might, when presented with a DNS name, generate an internal name that contained both the original DNS name and the equivalent X.500 name. Alternatively, it might only store the X.500 name. In the latter case, `gss_display_name` would most likely generate a printable X.500 name, rather than the original DNS name.

The process of authentication delivers to the context acceptor an internal name. Since this name has been authenticated by a single mechanism, it contains only a single name (even if the internal name presented by the context initiator to `gss_init_sec_context` had multiple components). Such names are termed internal mechanism names, or "MN"s and the names emitted by `gss_accept_sec_context` are always of this type. Since some applications may require MNs without wanting to incur the overhead of an authentication operation, a second function, `gss_canonicalize_name`, is provided to convert a general internal name into an MN.

Comparison of internal-form names may be accomplished via the `gss_compare_name` routine, which returns true if the two names being compared refer to the same entity. This removes the need for the application program to understand the syntaxes of the various printable names that a given GSS-API implementation may support. Since GSS-API assumes that all primitive names contained within a given internal name refer to the same entity, `gss_compare_name` can return true if the two names have at least one primitive name in common. If the implementation embodies knowledge of equivalence relationships between names taken from different namespaces, this knowledge may also allow successful comparison of internal names containing no overlapping primitive elements.

When used in large access control lists, the overhead of invoking `gss_import_name` and `gss_compare_name` on each name from the ACL may be prohibitive. As an alternative way of supporting this case, GSS-API defines a special form of the contiguous string name which may be compared directly (e.g. with `memcmp()`). Contiguous names suitable for comparison are generated by the `gss_export_name` routine, which requires an MN as input. Exported names may be re-imported by the `gss_import_name` routine, and the resulting internal name will also be an MN. The `gss_OID` constant `GSS_C_NT_EXPORT_NAME` identifies the "export name" type, and the value of this constant is given in Appendix A. Structurally, an exported name object consists of a header containing an OID identifying the mechanism that authenticated the name, and a trailer containing the name itself, where the syntax of the trailer is defined by the individual mechanism specification. The precise format of an export name is defined in the language-independent GSS-API specification [GSSAPI].

Note that the results obtained by using `gss_compare_name` will in general be different from those obtained by invoking `gss_canonicalize_name` and `gss_export_name`, and then comparing the exported names. The first series of operation determines whether two (unauthenticated) names identify the same principal; the second whether a particular mechanism would authenticate them as the same principal. These two operations will in general give the same results only for MNs.

The `gss_name_t` datatype should be implemented as a pointer type. To allow the compiler to aid the application programmer by performing type-checking, the use of `(void *)` is discouraged. A pointer to an implementation-defined type is the preferred choice.

Storage is allocated by routines that return `gss_name_t` values. A procedure, `gss_release_name`, is provided to free storage associated with an internal-form name.

3.2.6 Channel Bindings

GSS-API supports the use of user-specified tags to identify a given context to the peer application. These tags are intended to be used to identify the particular communications channel that carries the context. Channel bindings are communicated to the GSS-API using the following structure:

```

typedef struct gss_channel_bindings_struct {
    OM_uint32      initiator_addrtype;
    gss_buffer_desc initiator_address;
    OM_uint32      acceptor_addrtype;
    gss_buffer_desc acceptor_address;
    gss_buffer_desc application_data;
} *gss_channel_bindings_t;

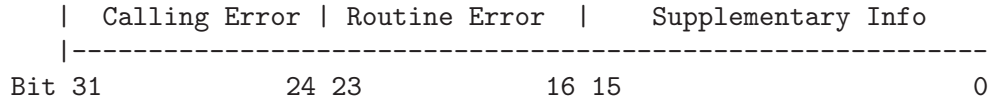
```

The `initiator_addrtype` and `acceptor_addrtype` fields denote the type of addresses contained in the `initiator_address` and `acceptor_address` buffers. The address type should be one of the following:

<code>GSS_C_AF_UNSPEC</code>	Unspecified address type
<code>GSS_C_AF_LOCAL</code>	Host-local address type
<code>GSS_C_AF_INET</code>	Internet address type (e.g. IP)
<code>GSS_C_AF_IMPLINK</code>	ARPAnet IMP address type
<code>GSS_C_AF_PUP</code>	pup protocols (eg BSP) address type
<code>GSS_C_AF_CHAOS</code>	MIT CHAOS protocol address type
<code>GSS_C_AF_NS</code>	XEROX NS address type
<code>GSS_C_AF_NBS</code>	nbs address type
<code>GSS_C_AF_ECMA</code>	ECMA address type
<code>GSS_C_AF_DATAKIT</code>	datakit protocols address type
<code>GSS_C_AF_CCITT</code>	CCITT protocols
<code>GSS_C_AF_SNA</code>	IBM SNA address type
<code>GSS_C_AF_DECnet</code>	DECnet address type
<code>GSS_C_AF_DLI</code>	Direct data link interface address type
<code>GSS_C_AF_LAT</code>	LAT address type
<code>GSS_C_AF_HYLINK</code>	NSC Hyperchannel address type
<code>GSS_C_AF_APPLETALK</code>	AppleTalk address type
<code>GSS_C_AF_BSC</code>	BISYNC 2780/3780 address type
<code>GSS_C_AF_DSS</code>	Distributed system services address type
<code>GSS_C_AF_OSI</code>	OSI TP4 address type
<code>GSS_C_AF_X25</code>	X.25
<code>GSS_C_AF_NULLADDR</code>	No address specified

Note that these symbols name address families rather than specific addressing formats. For address families that contain several alternative address forms, the `initiator_address` and `acceptor_address` fields must contain sufficient information to determine which address form is used. When not otherwise specified, addresses should be specified in network byte-order (that is, native byte-ordering for the address family).

Conceptually, the GSS-API concatenates the `initiator_addrtype`, `initiator_address`, `acceptor_addrtype`, `acceptor_address` and `application_data` to form an octet string. The mechanism calculates a MIC over this octet string, and binds the MIC to the context establishment token emitted by `gss_init_sec_context`. The same bindings are presented by the context acceptor to `gss_accept_sec_context`, and a MIC is calculated in the same way. The calculated MIC is compared with that found in the token, and if the MICs differ, `gss_accept_sec_context` will return a `GSS_S_BAD_BINDINGS` error, and the context will not be established. Some mechanisms may include the actual channel binding data in the



Hence if a GSS-API routine returns a GSS status code whose upper 16 bits contain a non-zero value, the call failed. If the calling error field is non-zero, the invoking application's call of the routine was erroneous. Calling errors are defined in table 3-1. If the routine error field is non-zero, the routine failed for one of the routine-specific reasons listed below in table 3-2. Whether or not the upper 16 bits indicate a failure or a success, the routine may indicate additional information by setting bits in the supplementary info field of the status code. The meaning of individual bits is listed below in table 3-3.

Table 3-1 Calling Errors

Name	Value in field	Meaning
GSS_S_CALL_INACCESSIBLE_READ	1	A required input parameter could not be read
GSS_S_CALL_INACCESSIBLE_WRITE	2	A required output parameter could not be written.
GSS_S_CALL_BAD_STRUCTURE	3	A parameter was malformed

Table 3-2 Routine Errors

Name	Value in field	Meaning	
GSS_S_BAD_MECH	1	An unsupported mechanism was requested	
GSS_S_BAD_NAME	2	An invalid name was supplied	
GSS_S_BAD_NAME_TYPE	3	A supplied name was of an unsupported type	
GSS_S_BAD_BINDINGS	4	Incorrect channel bindings were supplied	
GSS_S_BAD_STATUS	5	An invalid status code was supplied	
GSS_S_BAD_MIC	GSS_S_BAD_SIG	6	A token had an invalid MIC
GSS_S_NO_CRED	7	No credentials were supplied, or the credentials were unavailable or inaccessible.	
GSS_S_NO_CONTEXT	8	No context has been established	
GSS_S_DEFECTIVE_TOKEN	9	A token was invalid	
GSS_S_DEFECTIVE_CREDENTIAL	10	A credential was invalid	
GSS_S_CREDENTIALS_EXPIRED	11	The referenced credentials have expired	
GSS_S_CONTEXT_EXPIRED	12	The context has expired	

GSS_S_FAILURE	13	Miscellaneous failure (see text)
GSS_S_BAD_QOP	14	The quality-of-protection requested could not be provided
GSS_S_UNAUTHORIZED	15	The operation is forbidden by local security policy
GSS_S_UNAVAILABLE	16	The operation or option is unavailable
GSS_S_DUPLICATE_ELEMENT	17	The requested credential element already exists
GSS_S_NAME_NOT_MN	18	The provided name was not a mechanism name

Table 3-3 Supplementary Status Bits

Name	Bit Number	Meaning
-----	-----	-----
GSS_S_CONTINUE_NEEDED	0 (LSB)	Returned only by <code>gss_init_sec_context</code> or <code>gss_accept_sec_context</code> . The routine must be called again to complete its function. See routine documentation for detailed description
GSS_S_DUPLICATE_TOKEN	1	The token was a duplicate of an earlier token
GSS_S_OLD_TOKEN	2	The token's validity period has expired
GSS_S_UNSEQ_TOKEN	3	A later token has already been processed
GSS_S_GAP_TOKEN	4	An expected per-message token was not received

The routine documentation also uses the name `GSS_S_COMPLETE`, which is a zero value, to indicate an absence of any API errors or supplementary information bits.

All `GSS_S_XXX` symbols equate to complete `OM_uint32` status codes, rather than to bitfield values. For example, the actual value of the symbol `GSS_S_BAD_NAME_TYPE` (value 3 in the routine error field) is `3<<16`. The macros `GSS_CALLING_ERROR`, `GSS_ROUTINE_ERROR` and `GSS_SUPPLEMENTARY_INFO` are provided, each of which takes a GSS status code and removes all but the relevant field. For example, the value obtained by applying `GSS_ROUTINE_ERROR` to a status code removes the calling errors and supplementary info fields, leaving only the routine errors field. The values delivered by these macros may be directly compared with a `GSS_S_XXX` symbol of the appropriate type. The macro `GSS_ERROR` is also provided, which when applied to a GSS status code returns a non-zero value if the status code indicated a calling or routine error, and a zero value otherwise. All macros defined by GSS-API evaluate their argument(s) exactly once.

A GSS-API implementation may choose to signal calling errors in a platform-specific manner instead of, or in addition to the routine value; routine errors and supplementary info should be returned via major status values only.

The GSS major status code `GSS_S_FAILURE` is used to indicate that the underlying mechanism detected an error for which no specific GSS status code is defined. The mechanism-specific status code will provide more details about the error.

In addition to the explicit major status codes for each API function, the code `GSS_S_FAILURE` may be returned by any routine, indicating an implementation-specific or mechanism-specific error condition, further details of which are reported via the `minor_status` parameter.

3.4.2 Mechanism-specific status codes

GSS-API routines return a `minor_status` parameter, which is used to indicate specialized errors from the underlying security mechanism. This parameter may contain a single mechanism-specific error, indicated by a `OM_uint32` value.

The `minor_status` parameter will always be set by a GSS-API routine, even if it returns a calling error or one of the generic API errors indicated above as fatal, although most other output parameters may remain unset in such cases. However, output parameters that are expected to return pointers to storage allocated by a routine must always be set by the routine, even in the event of an error, although in such cases the GSS-API routine may elect to set the returned parameter value to `NULL` to indicate that no storage was actually allocated. Any length field associated with such pointers (as in a `gss_buffer_desc` structure) should also be set to zero in such cases.

3.5 Credential Management

GSS-API Credential-management Routines

Routine -----	Function -----
<code>gss_acquire_cred</code>	Assume a global identity; Obtain a GSS-API credential handle for pre-existing credentials.
<code>gss_add_cred</code>	Construct credentials incrementally.
<code>gss_inquire_cred</code>	Obtain information about a credential.
<code>gss_inquire_cred_by_mech</code>	Obtain per-mechanism information about a credential.
<code>gss_release_cred</code>	Discard a credential handle.

`OM_uint32 gss_acquire_cred (OM_uint32 * minor_status, const [Function]
 gss_name_t desired_name, OM_uint32 time_req, const gss_OID_set
 desired_mechs, gss_cred_usage_t cred_usage, gss_cred_id_t *
 output_cred_handle, gss_OID_set * actual_mechs, OM_uint32 *
 time_rec)
 minor_status: (integer, modify) Mechanism specific status code.`

desired_name: (gss_name_t, read) Name of principal whose credential should be acquired.

time_req: (Integer, read, optional) Number of seconds that credentials should remain valid. Specify GSS_C_INDEFINITE to request that the credentials have the maximum permitted lifetime.

desired_mechs: (Set of Object IDs, read, optional) Set of underlying security mechanisms that may be used. GSS_C_NO_OID_SET may be used to obtain an implementation-specific default.

cred_usage: (gss_cred_usage_t, read) GSS_C_BOTH - Credentials may be used either to initiate or accept security contexts. GSS_C_INITIATE - Credentials will only be used to initiate security contexts. GSS_C_ACCEPT - Credentials will only be used to accept security contexts.

output_cred_handle: (gss_cred_id_t, modify) The returned credential handle. Resources associated with this credential handle must be released by the application after use with a call to gss_release_cred().

actual_mechs: (Set of Object IDs, modify, optional) The set of mechanisms for which the credential is valid. Storage associated with the returned OID-set must be released by the application after use with a call to gss_release_oid_set(). Specify NULL if not required.

time_rec: (Integer, modify, optional) Actual number of seconds for which the returned credentials will remain valid. If the implementation does not support expiration of credentials, the value GSS_C_INDEFINITE will be returned. Specify NULL if not required.

Allows an application to acquire a handle for a pre-existing credential by name. GSS-API implementations must impose a local access-control policy on callers of this routine to prevent unauthorized callers from acquiring credentials to which they are not entitled. This routine is not intended to provide a "login to the network" function, as such a function would involve the creation of new credentials rather than merely acquiring a handle to existing credentials. Such functions, if required, should be defined in implementation-specific extensions to the API.

If *desired_name* is GSS_C_NO_NAME, the call is interpreted as a request for a credential handle that will invoke default behavior when passed to gss_init_sec_context() (if *cred_usage* is GSS_C_INITIATE or GSS_C_BOTH) or gss_accept_sec_context() (if *cred_usage* is GSS_C_ACCEPT or GSS_C_BOTH).

Mechanisms should honor the *desired_mechs* parameter, and return a credential that is suitable to use only with the requested mechanisms. An exception to this is the case where one underlying credential element can be shared by multiple mechanisms; in this case it is permissible for an implementation to indicate all mechanisms with which the credential element may be used. If *desired_mechs* is an empty set, behavior is undefined.

This routine is expected to be used primarily by context acceptors, since implementations are likely to provide mechanism-specific ways of obtaining GSS-API initiator credentials from the system login process. Some implementations may therefore not support the acquisition of GSS_C_INITIATE or GSS_C_BOTH credentials via gss_acquire_cred for any name other than GSS_C_NO_NAME, or a name produced

by applying either `gss_inquire_cred` to a valid credential, or `gss_inquire_context` to an active context.

If credential acquisition is time-consuming for a mechanism, the mechanism may choose to delay the actual acquisition until the credential is required (e.g. by `gss_init_sec_context` or `gss_accept_sec_context`). Such mechanism-specific implementation decisions should be invisible to the calling application; thus a call of `gss_inquire_cred` immediately following the call of `gss_acquire_cred` must return valid credential data, and may therefore incur the overhead of a deferred credential acquisition.

Return value:

`GSS_S_COMPLETE`: Successful completion.

`GSS_S_BAD_MECH`: Unavailable mechanism requested.

`GSS_S_BAD_NAME_TYPE`: Type contained within `desired_name` parameter is not supported.

`GSS_S_BAD_NAME`: Value supplied for `desired_name` parameter is ill formed.

`GSS_S_CREDENTIALS_EXPIRED`: The credentials could not be acquired Because they have expired.

`GSS_S_NO_CRED`: No credentials were found for the specified name.

```
OM_uint32 gss_add_cred (OM_uint32 * minor_status, const          [Function]
    gss_cred_id_t input_cred_handle, const gss_name_t desired_name, const
    gss_OID desired_mech, gss_cred_usage_t cred_usage, OM_uint32
    initiator_time_req, OM_uint32 acceptor_time_req, gss_cred_id_t *
    output_cred_handle, gss_OID_set * actual_mechs, OM_uint32 *
    initiator_time_rec, OM_uint32 * acceptor_time_rec)
```

minor_status: (integer, modify) Mechanism specific status code.

input_cred_handle: (gss_cred_id_t, read, optional) The credential to which a credential-element will be added. If `GSS_C_NO_CREDENTIAL` is specified, the routine will compose the new credential based on default behavior (see text). Note that, while the credential-handle is not modified by `gss_add_cred()`, the underlying credential will be modified if `output_cred_handle` is `NULL`.

desired_name: (gss_name_t, read.) Name of principal whose credential should be acquired.

desired_mech: (Object ID, read) Underlying security mechanism with which the credential may be used.

cred_usage: (gss_cred_usage_t, read) `GSS_C_BOTH` - Credential may be used either to initiate or accept security contexts. `GSS_C_INITIATE` - Credential will only be used to initiate security contexts. `GSS_C_ACCEPT` - Credential will only be used to accept security contexts.

initiator_time_req: (Integer, read, optional) number of seconds that the credential should remain valid for initiating security contexts. This argument is ignored if the composed credentials are of type `GSS_C_ACCEPT`. Specify `GSS_C_INDEFINITE` to request that the credentials have the maximum permitted initiator lifetime.

acceptor_time_req: (Integer, read, optional) number of seconds that the credential should remain valid for accepting security contexts. This argument is ignored if the composed credentials are of type GSS_C_INITIATE. Specify GSS_C_INDEFINITE to request that the credentials have the maximum permitted initiator lifetime.

output_cred_handle: (gss_cred_id_t, modify, optional) The returned credential handle, containing the new credential-element and all the credential-elements from *input_cred_handle*. If a valid pointer to a gss_cred_id_t is supplied for this parameter, gss_add_cred creates a new credential handle containing all credential-elements from the *input_cred_handle* and the newly acquired credential-element; if NULL is specified for this parameter, the newly acquired credential-element will be added to the credential identified by *input_cred_handle*. The resources associated with any credential handle returned via this parameter must be released by the application after use with a call to gss_release_cred().

actual_mechs: (Set of Object IDs, modify, optional) The complete set of mechanisms for which the new credential is valid. Storage for the returned OID-set must be freed by the application after use with a call to gss_release_oid_set(). Specify NULL if not required.

initiator_time_rec: (Integer, modify, optional) Actual number of seconds for which the returned credentials will remain valid for initiating contexts using the specified mechanism. If the implementation or mechanism does not support expiration of credentials, the value GSS_C_INDEFINITE will be returned. Specify NULL if not required

acceptor_time_rec: (Integer, modify, optional) Actual number of seconds for which the returned credentials will remain valid for accepting security contexts using the specified mechanism. If the implementation or mechanism does not support expiration of credentials, the value GSS_C_INDEFINITE will be returned. Specify NULL if not required

Adds a credential-element to a credential. The credential-element is identified by the name of the principal to which it refers. GSS-API implementations must impose a local access-control policy on callers of this routine to prevent unauthorized callers from acquiring credential-elements to which they are not entitled. This routine is not intended to provide a "login to the network" function, as such a function would involve the creation of new mechanism-specific authentication data, rather than merely acquiring a GSS-API handle to existing data. Such functions, if required, should be defined in implementation-specific extensions to the API.

If *desired_name* is GSS_C_NO_NAME, the call is interpreted as a request to add a credential element that will invoke default behavior when passed to gss_init_sec_context() (if *cred_usage* is GSS_C_INITIATE or GSS_C_BOTH) or gss_accept_sec_context() (if *cred_usage* is GSS_C_ACCEPT or GSS_C_BOTH).

This routine is expected to be used primarily by context acceptors, since implementations are likely to provide mechanism-specific ways of obtaining GSS-API initiator credentials from the system login process. Some implementations may therefore not support the acquisition of GSS_C_INITIATE or GSS_C_BOTH credentials via gss_acquire_cred for any name other than GSS_C_NO_NAME, or a name produced

by applying either `gss_inquire_cred` to a valid credential, or `gss_inquire_context` to an active context.

If credential acquisition is time-consuming for a mechanism, the mechanism may choose to delay the actual acquisition until the credential is required (e.g. by `gss_init_sec_context` or `gss_accept_sec_context`). Such mechanism-specific implementation decisions should be invisible to the calling application; thus a call of `gss_inquire_cred` immediately following the call of `gss_add_cred` must return valid credential data, and may therefore incur the overhead of a deferred credential acquisition.

This routine can be used to either compose a new credential containing all credential-elements of the original in addition to the newly-acquire credential-element, or to add the new credential- element to an existing credential. If `NULL` is specified for the `output_cred_handle` parameter argument, the new credential-element will be added to the credential identified by `input_cred_handle`; if a valid pointer is specified for the `output_cred_handle` parameter, a new credential handle will be created.

If `GSS_C_NO_CREDENTIAL` is specified as the `input_cred_handle`, `gss_add_cred` will compose a credential (and set the `output_cred_handle` parameter accordingly) based on default behavior. That is, the call will have the same effect as if the application had first made a call to `gss_acquire_cred()`, specifying the same usage and passing `GSS_C_NO_NAME` as the `desired_name` parameter to obtain an explicit credential handle embodying default behavior, passed this credential handle to `gss_add_cred()`, and finally called `gss_release_cred()` on the first credential handle.

If `GSS_C_NO_CREDENTIAL` is specified as the `input_cred_handle` parameter, a non-`NULL` `output_cred_handle` must be supplied.

Return value:

`GSS_S_COMPLETE`: Successful completion.

`GSS_S_BAD_MECH`: Unavailable mechanism requested.

`GSS_S_BAD_NAME_TYPE`: Type contained within `desired_name` parameter is not supported.

`GSS_S_BAD_NAME`: Value supplied for `desired_name` parameter is ill-formed.

`GSS_S_DUPLICATE_ELEMENT`: The credential already contains an element for the requested mechanism with overlapping usage and validity period.

`GSS_S_CREDENTIALS_EXPIRED`: The required credentials could not be added because they have expired.

`GSS_S_NO_CRED`: No credentials were found for the specified name.

```
OM_uint32 gss_inquire_cred (OM_uint32 * minor_status, const [Function]
    gss_cred_id_t cred_handle, gss_name_t * name, OM_uint32 * lifetime,
    gss_cred_usage_t * cred_usage, gss_OID_set * mechanisms)
minor_status: (integer, modify) Mechanism specific status code.
```

`cred_handle`: (`gss_cred_id_t`, read) A handle that refers to the target credential. Specify `GSS_C_NO_CREDENTIAL` to inquire about the default initiator principal.

`name`: (`gss_name_t`, modify, optional) The name whose identity the credential asserts. Storage associated with this name should be freed by the application after use with a call to `gss_release_name()`. Specify `NULL` if not required.

lifetime: (Integer, modify, optional) The number of seconds for which the credential will remain valid. If the credential has expired, this parameter will be set to zero. If the implementation does not support credential expiration, the value GSS_C_INDEFINITE will be returned. Specify NULL if not required.

cred_usage: (gss_cred_usage_t, modify, optional) How the credential may be used. One of the following: GSS_C_INITIATE, GSS_C_ACCEPT, GSS_C_BOTH. Specify NULL if not required.

mechanisms: (gss_OID_set, modify, optional) Set of mechanisms supported by the credential. Storage associated with this OID set must be freed by the application after use with a call to gss_release_oid_set(). Specify NULL if not required.

Obtains information about a credential.

Return value:

GSS_S_COMPLETE: Successful completion.

GSS_S_NO_CRED: The referenced credentials could not be accessed.

GSS_S_DEFECTIVE_CREDENTIAL: The referenced credentials were invalid.

GSS_S_CREDENTIALS_EXPIRED: The referenced credentials have expired. If the lifetime parameter was not passed as NULL, it will be set to 0.

```
OM_uint32 gss_inquire_cred_by_mech (OM_uint32 *                               [Function]
    minor_status, const gss_cred_id_t cred_handle, const gss_OID mech_type,
    gss_name_t * name, OM_uint32 * initiator_lifetime, OM_uint32 *
    acceptor_lifetime, gss_cred_usage_t * cred_usage)
```

minor_status: (Integer, modify) Mechanism specific status code.

cred_handle: (gss_cred_id_t, read) A handle that refers to the target credential. Specify GSS_C_NO_CREDENTIAL to inquire about the default initiator principal.

mech_type: (gss_OID, read) The mechanism for which information should be returned.

name: (gss_name_t, modify, optional) The name whose identity the credential asserts. Storage associated with this name must be freed by the application after use with a call to gss_release_name(). Specify NULL if not required.

initiator_lifetime: (Integer, modify, optional) The number of seconds for which the credential will remain capable of initiating security contexts under the specified mechanism. If the credential can no longer be used to initiate contexts, or if the credential usage for this mechanism is GSS_C_ACCEPT, this parameter will be set to zero. If the implementation does not support expiration of initiator credentials, the value GSS_C_INDEFINITE will be returned. Specify NULL if not required.

acceptor_lifetime: (Integer, modify, optional) The number of seconds for which the credential will remain capable of accepting security contexts under the specified mechanism. If the credential can no longer be used to accept contexts, or if the credential usage for this mechanism is GSS_C_INITIATE, this parameter will be set to zero. If the implementation does not support expiration of acceptor credentials, the value GSS_C_INDEFINITE will be returned. Specify NULL if not required.

cred_usage: (gss_cred_usage_t, modify, optional) How the credential may be used with the specified mechanism. One of the following: GSS_C_INITIATE, GSS_C_ACCEPT, GSS_C_BOTH. Specify NULL if not required.

Obtains per-mechanism information about a credential.

Return value:

GSS_S_COMPLETE: Successful completion.

GSS_S_NO_CRED: The referenced credentials could not be accessed.

GSS_S_DEFECTIVE_CREDENTIAL: The referenced credentials were invalid.

GSS_S_CREDENTIALS_EXPIRED: The referenced credentials have expired. If the lifetime parameter was not passed as NULL, it will be set to 0.

OM_uint32 gss_release_cred (OM_uint32 * minor_status, [Function]
 gss_cred_id_t * cred_handle)

minor_status: (Integer, modify) Mechanism specific status code.

cred_handle: (gss_cred_id_t, modify, optional) Opaque handle identifying credential to be released. If GSS_C_NO_CREDENTIAL is supplied, the routine will complete successfully, but will do nothing.

Informs GSS-API that the specified credential handle is no longer required by the application, and frees associated resources. The cred_handle is set to GSS_C_NO_CREDENTIAL on successful completion of this call.

Return value:

GSS_S_COMPLETE: Successful completion.

GSS_S_NO_CRED: Credentials could not be accessed.

3.6 Context-Level Routines

GSS-API Context-Level Routines

Routine	Function
-----	-----
gss_init_sec_context	Initiate a security context with a peer application.
gss_accept_sec_context	Accept a security context initiated by a peer application.
gss_delete_sec_context	Discard a security context.
gss_process_context_token	Process a token on a security context from a peer application.
gss_context_time	Determine for how long a context will remain valid.
gss_inquire_context	Obtain information about a security context.
gss_wrap_size_limit	Determine token-size limit for gss_wrap on a context.
gss_export_sec_context	Transfer a security context to another process.
gss_import_sec_context	Import a transferred context.

```

OM_uint32 gss_init_sec_context (OM_uint32 *minor_status,          [Function]
    const gss_cred_id_t initiator_cred_handle, gss_ctx_id_t *
    context_handle, const gss_name_t target_name, const gss_OID
    mech_type, OM_uint32 req_flags, OM_uint32 time_req, const
    gss_channel_bindings_t input_chan_bindings, const gss_buffer_t
    input_token, gss_OID * actual_mech_type, gss_buffer_t output_token,
    OM_uint32 * ret_flags, OM_uint32 * time_rec)

```

minor_status: (integer, modify) Mechanism specific status code.

initiator_cred_handle: (gss_cred_id_t, read, optional) Handle for credentials claimed. Supply GSS_C_NO_CREDENTIAL to act as a default initiator principal. If no default initiator is defined, the function will return GSS_S_NO_CRED.

context_handle: (gss_ctx_id_t, read/modify) Context handle for new context. Supply GSS_C_NO_CONTEXT for first call; use value returned by first call in continuation calls. Resources associated with this context-handle must be released by the application after use with a call to `gss_delete_sec_context()`.

target_name: (gss_name_t, read) Name of target.

mech_type: (OID, read, optional) Object ID of desired mechanism. Supply GSS_C_NO_OID to obtain an implementation specific default.

req_flags: (bit-mask, read) Contains various independent flags, each of which requests that the context support a specific service option. Symbolic names are provided for each flag, and the symbolic names corresponding to the required flags should be logically-ORed together to form the bit-mask value. See below for the flags.

time_req: (Integer, read, optional) Desired number of seconds for which context should remain valid. Supply 0 to request a default validity period.

input_chan_bindings: (channel bindings, read, optional) Application-specified bindings. Allows application to securely bind channel identification information to the security context. Specify GSS_C_NO_CHANNEL_BINDINGS if channel bindings are not used.

input_token: (buffer, opaque, read, optional) Token received from peer application. Supply GSS_C_NO_BUFFER, or a pointer to a buffer containing the value GSS_C_EMPTY_BUFFER on initial call.

actual_mech_type: (OID, modify, optional) Actual mechanism used. The OID returned via this parameter will be a pointer to static storage that should be treated as read-only; In particular the application should not attempt to free it. Specify NULL if not required.

output_token: (buffer, opaque, modify) Token to be sent to peer application. If the length field of the returned buffer is zero, no token need be sent to the peer application. Storage associated with this buffer must be freed by the application after use with a call to `gss_release_buffer()`.

ret_flags: (bit-mask, modify, optional) Contains various independent flags, each of which indicates that the context supports a specific service option. Specify NULL if not required. Symbolic names are provided for each flag, and the symbolic names corresponding to the required flags should be logically-ANDed with the `ret_flags` value to test whether a given option is supported by the context. See below for the flags.

time_rec: (Integer, modify, optional) Number of seconds for which the context will remain valid. If the implementation does not support context expiration, the value `GSS_C_INDEFINITE` will be returned. Specify `NULL` if not required.

Initiates the establishment of a security context between the application and a remote peer. Initially, the `input_token` parameter should be specified either as `GSS_C_NO_BUFFER`, or as a pointer to a `gss_buffer_desc` object whose length field contains the value zero. The routine may return a `output_token` which should be transferred to the peer application, where the peer application will present it to `gss_accept_sec_context`. If no token need be sent, `gss_init_sec_context` will indicate this by setting the length field of the `output_token` argument to zero. To complete the context establishment, one or more reply tokens may be required from the peer application; if so, `gss_init_sec_context` will return a status containing the supplementary information bit `GSS_S_CONTINUE_NEEDED`. In this case, `gss_init_sec_context` should be called again when the reply token is received from the peer application, passing the reply token to `gss_init_sec_context` via the `input_token` parameters.

Portable applications should be constructed to use the token length and return status to determine whether a token needs to be sent or waited for. Thus a typical portable caller should always invoke `gss_init_sec_context` within a loop:

```
int context_established = 0;
gss_ctx_id_t context_hdl = GSS_C_NO_CONTEXT;
...
input_token->length = 0;

while (!context_established) {
    maj_stat = gss_init_sec_context(&min_stat,
                                   cred_hdl,
                                   &context_hdl,
                                   target_name,
                                   desired_mech,
                                   desired_services,
                                   desired_time,
                                   input_bindings,
                                   input_token,
                                   &actual_mech,
                                   output_token,
                                   &actual_services,
                                   &actual_time);

    if (GSS_ERROR(maj_stat)) {
        report_error(maj_stat, min_stat);
    };

    if (output_token->length != 0) {
        send_token_to_peer(output_token);
        gss_release_buffer(&min_stat, output_token)
    };
};
```

```

if (GSS_ERROR(maj_stat)) {

    if (context_hdl != GSS_C_NO_CONTEXT)
        gss_delete_sec_context(&min_stat,
                               &context_hdl,
                               GSS_C_NO_BUFFER);

    break;
};

if (maj_stat & GSS_S_CONTINUE_NEEDED) {
    receive_token_from_peer(input_token);
} else {
    context_established = 1;
};
};

```

Whenever the routine returns a major status that includes the value `GSS_S_CONTINUE_NEEDED`, the context is not fully established and the following restrictions apply to the output parameters:

- The value returned via the `time_rec` parameter is undefined unless the accompanying `ret_flags` parameter contains the bit `GSS_C_PROT_READY_FLAG`, indicating that per-message services may be applied in advance of a successful completion status, the value returned via the `actual_mech_type` parameter is undefined until the routine returns a major status value of `GSS_S_COMPLETE`.
- The values of the `GSS_C_DELEG_FLAG`, `GSS_C_MUTUAL_FLAG`, `GSS_C_REPLAY_FLAG`, `GSS_C_SEQUENCE_FLAG`, `GSS_C_CONF_FLAG`, `GSS_C_INTEG_FLAG` and `GSS_C_ANON_FLAG` bits returned via the `ret_flags` parameter should contain the values that the implementation expects would be valid if context establishment were to succeed. In particular, if the application has requested a service such as delegation or anonymous authentication via the `req_flags` argument, and such a service is unavailable from the underlying mechanism, `gss_init_sec_context` should generate a token that will not provide the service, and indicate via the `ret_flags` argument that the service will not be supported. The application may choose to abort the context establishment by calling `gss_delete_sec_context` (if it cannot continue in the absence of the service), or it may choose to transmit the token and continue context establishment (if the service was merely desired but not mandatory).
- The values of the `GSS_C_PROT_READY_FLAG` and `GSS_C_TRANS_FLAG` bits within `ret_flags` should indicate the actual state at the time `gss_init_sec_context` returns, whether or not the context is fully established.
- GSS-API implementations that support per-message protection are encouraged to set the `GSS_C_PROT_READY_FLAG` in the final `ret_flags` returned to a caller (i.e. when accompanied by a `GSS_S_COMPLETE` status code). However, applications should not rely on this behavior as the flag was not defined in Version 1 of the GSS-API. Instead, applications should determine what per-message services are available after a successful context establishment according to the `GSS_C_INTEG_FLAG` and `GSS_C_CONF_FLAG` values.

- All other bits within the `ret_flags` argument should be set to zero.

If the initial call of `gss_init_sec_context()` fails, the implementation should not create a context object, and should leave the value of the `context_handle` parameter set to `GSS_C_NO_CONTEXT` to indicate this. In the event of a failure on a subsequent call, the implementation is permitted to delete the "half-built" security context (in which case it should set the `context_handle` parameter to `GSS_C_NO_CONTEXT`), but the preferred behavior is to leave the security context untouched for the application to delete (using `gss_delete_sec_context`).

During context establishment, the informational status bits `GSS_S_OLD_TOKEN` and `GSS_S_DUPLICATE_TOKEN` indicate fatal errors, and GSS-API mechanisms should always return them in association with a routine error of `GSS_S_FAILURE`. This requirement for pairing did not exist in version 1 of the GSS-API specification, so applications that wish to run over version 1 implementations must special-case these codes.

The `req_flags` values:

`GSS_C_DELEG_FLAG`

- True - Delegate credentials to remote peer.
- False - Don't delegate.

`GSS_C_MUTUAL_FLAG`

- True - Request that remote peer authenticate itself.
- False - Authenticate self to remote peer only.

`GSS_C_REPLAY_FLAG`

- True - Enable replay detection for messages protected with `gss_wrap` or `gss_get_mic`.
- False - Don't attempt to detect replayed messages.

`GSS_C_SEQUENCE_FLAG`

- True - Enable detection of out-of-sequence protected messages.
- False - Don't attempt to detect out-of-sequence messages.

`GSS_C_CONF_FLAG`

- True - Request that confidentiality service be made available (via `gss_wrap`).
- False - No per-message confidentiality service is required.

`GSS_C_INTEG_FLAG`

- True - Request that integrity service be made available (via `gss_wrap` or `gss_get_mic`).
- False - No per-message integrity service is required.

`GSS_C_ANON_FLAG`

- True - Do not reveal the initiator's identity to the acceptor.
- False - Authenticate normally.

The `ret_flags` values:

GSS_C_DELEG_FLAG

- True - Credentials were delegated to the remote peer.
- False - No credentials were delegated.

GSS_C_MUTUAL_FLAG

- True - The remote peer has authenticated itself.
- False - Remote peer has not authenticated itself.

GSS_C_REPLAY_FLAG

- True - replay of protected messages will be detected.
- False - replayed messages will not be detected.

GSS_C_SEQUENCE_FLAG

- True - out-of-sequence protected messages will be detected.
- False - out-of-sequence messages will not be detected.

GSS_C_CONF_FLAG

- True - Confidentiality service may be invoked by calling `gss_wrap` routine.
- False - No confidentiality service (via `gss_wrap`) available. `gss_wrap` will provide message encapsulation, data-origin authentication and integrity services only.

GSS_C_INTEG_FLAG

- True - Integrity service may be invoked by calling either `gss_get_mic` or `gss_wrap` routines.
- False - Per-message integrity service unavailable.

GSS_C_ANON_FLAG

- True - The initiator's identity has not been revealed, and will not be revealed if any emitted token is passed to the acceptor.
- False - The initiator's identity has been or will be authenticated normally.

GSS_C_PROT_READY_FLAG

- True - Protection services (as specified by the states of the `GSS_C_CONF_FLAG` and `GSS_C_INTEG_FLAG`) are available for use if the accompanying major status return value is either `GSS_S_COMPLETE` or `GSS_S_CONTINUE_NEEDED`.
- False - Protection services (as specified by the states of the `GSS_C_CONF_FLAG` and `GSS_C_INTEG_FLAG`) are available only if the accompanying major status return value is `GSS_S_COMPLETE`.

GSS_C_TRANS_FLAG

- True - The resultant security context may be transferred to other processes via a call to `gss_export_sec_context()`.

- False - The security context is not transferable.

All other bits should be set to zero.

Return value:

GSS_S_COMPLETE: Successful completion.

GSS_S_CONTINUE_NEEDED: Indicates that a token from the peer application is required to complete the context, and that `gss_init_sec_context` must be called again with that token.

GSS_S_DEFECTIVE_TOKEN: Indicates that consistency checks performed on the `input_token` failed.

GSS_S_DEFECTIVE_CREDENTIAL: Indicates that consistency checks performed on the credential failed.

GSS_S_NO_CRED: The supplied credentials were not valid for context initiation, or the credential handle did not reference any credentials.

GSS_S_CREDENTIALS_EXPIRED: The referenced credentials have expired.

GSS_S_BAD_BINDINGS: The `input_token` contains different channel bindings to those specified via the `input_chan_bindings` parameter.

GSS_S_BAD_SIG: The `input_token` contains an invalid MIC, or a MIC that could not be verified.

GSS_S_OLD_TOKEN: The `input_token` was too old. This is a fatal error during context establishment.

GSS_S_DUPLICATE_TOKEN: The `input_token` is valid, but is a duplicate of a token already processed. This is a fatal error during context establishment.

GSS_S_NO_CONTEXT: Indicates that the supplied context handle did not refer to a valid context.

GSS_S_BAD_NAME_TYPE: The provided `target_name` parameter contained an invalid or unsupported type of name.

GSS_S_BAD_NAME: The provided `target_name` parameter was ill-formed.

GSS_S_BAD_MECH: The specified mechanism is not supported by the provided credential, or is unrecognized by the implementation.

```
OM_uint32 gss_accept_sec_context (OM_uint32 *minor_status,      [Function]
    gss_ctx_id_t *context_handle, const gss_cred_id_t
    acceptor_cred_handle, const gss_buffer_t input_token_buffer, const
    gss_channel_bindings_t input_chan_bindings, gss_name_t *src_name,
    gss_OID *mech_type, gss_buffer_t output_token, OM_uint32 *ret_flags,
    OM_uint32 *time_rec, gss_cred_id_t *delegated_cred_handle)
minor_status: (Integer, modify) Mechanism specific status code.
```

context_handle: (gss_ctx_id_t, read/modify) Context handle for new context. Supply GSS_C_NO_CONTEXT for first call; use value returned in subsequent calls. Once `gss_accept_sec_context()` has returned a value via this parameter, resources have been assigned to the corresponding context, and must be freed by the application after use with a call to `gss_delete_sec_context()`.

acceptor_cred_handle: (gss_cred_id_t, read) Credential handle claimed by context acceptor. Specify GSS_C_NO_CREDENTIAL to accept the context as a default principal. If GSS_C_NO_CREDENTIAL is specified, but no default acceptor principal is defined, GSS_S_NO_CRED will be returned.

input_token_buffer: (buffer, opaque, read) Token obtained from remote application.

input_chan_bindings: (channel bindings, read, optional) Application- specified bindings. Allows application to securely bind channel identification information to the security context. If channel bindings are not used, specify GSS_C_NO_CHANNEL_BINDINGS.

src_name: (gss_name_t, modify, optional) Authenticated name of context initiator. After use, this name should be deallocated by passing it to gss_release_name(). If not required, specify NULL.

mech_type: (Object ID, modify, optional) Security mechanism used. The returned OID value will be a pointer into static storage, and should be treated as read-only by the caller (in particular, it does not need to be freed). If not required, specify NULL.

output_token: (buffer, opaque, modify) Token to be passed to peer application. If the length field of the returned token buffer is 0, then no token need be passed to the peer application. If a non- zero length field is returned, the associated storage must be freed after use by the application with a call to gss_release_buffer().

ret_flags: (bit-mask, modify, optional) Contains various independent flags, each of which indicates that the context supports a specific service option. If not needed, specify NULL. Symbolic names are provided for each flag, and the symbolic names corresponding to the required flags should be logically-ANDed with the ret_flags value to test whether a given option is supported by the context. See below for the flags.

time_rec: (Integer, modify, optional) Number of seconds for which the context will remain valid. Specify NULL if not required.

delegated_cred_handle: (gss_cred_id_t, modify, optional credential) Handle for credentials received from context initiator. Only valid if deleg_flag in ret_flags is true, in which case an explicit credential handle (i.e. not GSS_C_NO_CREDENTIAL) will be returned; if deleg_flag is false, gss_accept_context() will set this parameter to GSS_C_NO_CREDENTIAL. If a credential handle is returned, the associated resources must be released by the application after use with a call to gss_release_cred(). Specify NULL if not required.

Allows a remotely initiated security context between the application and a remote peer to be established. The routine may return a output_token which should be transferred to the peer application, where the peer application will present it to gss_init_sec_context. If no token need be sent, gss_accept_sec_context will indicate this by setting the length field of the output_token argument to zero. To complete the context establishment, one or more reply tokens may be required from the peer application; if so, gss_accept_sec_context will return a status flag of GSS_S_CONTINUE_NEEDED, in which case it should be called again when the reply token is received from the peer application, passing the token to gss_accept_sec_context via the input_token parameters.

Portable applications should be constructed to use the token length and return status to determine whether a token needs to be sent or waited for. Thus a typical portable caller should always invoke `gss_accept_sec_context` within a loop:

```

gss_ctx_id_t context_hdl = GSS_C_NO_CONTEXT;

do {
    receive_token_from_peer(input_token);
    maj_stat = gss_accept_sec_context(&min_stat,
                                     &context_hdl,
                                     cred_hdl,
                                     input_token,
                                     input_bindings,
                                     &client_name,
                                     &mech_type,
                                     output_token,
                                     &ret_flags,
                                     &time_rec,
                                     &deleg_cred);

    if (GSS_ERROR(maj_stat)) {
        report_error(maj_stat, min_stat);
    };
    if (output_token->length != 0) {
        send_token_to_peer(output_token);

        gss_release_buffer(&min_stat, output_token);
    };
    if (GSS_ERROR(maj_stat)) {
        if (context_hdl != GSS_C_NO_CONTEXT)
            gss_delete_sec_context(&min_stat,
                                   &context_hdl,
                                   GSS_C_NO_BUFFER);

        break;
    };
} while (maj_stat & GSS_S_CONTINUE_NEEDED);

```

Whenever the routine returns a major status that includes the value `GSS_S_CONTINUE_NEEDED`, the context is not fully established and the following restrictions apply to the output parameters:

The value returned via the `time_rec` parameter is undefined Unless the accompanying `ret_flags` parameter contains the bit `GSS_C_PROT_READY_FLAG`, indicating that per-message services may be applied in advance of a successful completion status, the value returned via the `mech_type` parameter may be undefined until the routine returns a major status value of `GSS_S_COMPLETE`.

The values of the `GSS_C_DELEG_FLAG`, `GSS_C_MUTUAL_FLAG`, `GSS_C_REPLAY_FLAG`, `GSS_C_SEQUENCE_FLAG`, `GSS_C_CONF_FLAG`, `GSS_C_INTEG_FLAG` and `GSS_C_ANON_FLAG` bits returned via the `ret_flags` parameter should contain the

values that the implementation expects would be valid if context establishment were to succeed.

The values of the `GSS_C_PROT_READY_FLAG` and `GSS_C_TRANS_FLAG` bits within `ret_flags` should indicate the actual state at the time `gss_accept_sec_context` returns, whether or not the context is fully established.

Although this requires that GSS-API implementations set the `GSS_C_PROT_READY_FLAG` in the final `ret_flags` returned to a caller (i.e. when accompanied by a `GSS_S_COMPLETE` status code), applications should not rely on this behavior as the flag was not defined in Version 1 of the GSS-API. Instead, applications should be prepared to use per-message services after a successful context establishment, according to the `GSS_C_INTEG_FLAG` and `GSS_C_CONF_FLAG` values.

All other bits within the `ret_flags` argument should be set to zero. While the routine returns `GSS_S_CONTINUE_NEEDED`, the values returned via the `ret_flags` argument indicate the services that the implementation expects to be available from the established context.

If the initial call of `gss_accept_sec_context()` fails, the implementation should not create a context object, and should leave the value of the `context_handle` parameter set to `GSS_C_NO_CONTEXT` to indicate this. In the event of a failure on a subsequent call, the implementation is permitted to delete the "half-built" security context (in which case it should set the `context_handle` parameter to `GSS_C_NO_CONTEXT`), but the preferred behavior is to leave the security context (and the `context_handle` parameter) untouched for the application to delete (using `gss_delete_sec_context`).

During context establishment, the informational status bits `GSS_S_OLD_TOKEN` and `GSS_S_DUPLICATE_TOKEN` indicate fatal errors, and GSS-API mechanisms should always return them in association with a routine error of `GSS_S_FAILURE`. This requirement for pairing did not exist in version 1 of the GSS-API specification, so applications that wish to run over version 1 implementations must special-case these codes.

The `ret_flags` values:

`GSS_C_DELEG_FLAG`

- True - Delegated credentials are available via the `delegated_cred_handle` parameter.
- False - No credentials were delegated.

`GSS_C_MUTUAL_FLAG`

- True - Remote peer asked for mutual authentication.
- False - Remote peer did not ask for mutual authentication.

`GSS_C_REPLAY_FLAG`

- True - replay of protected messages will be detected.
- False - replayed messages will not be detected.

`GSS_C_SEQUENCE_FLAG`

- True - out-of-sequence protected messages will be detected.
- False - out-of-sequence messages will not be detected.

GSS_C_CONF_FLAG

- True - Confidentiality service may be invoked by calling the `gss_wrap` routine.
- False - No confidentiality service (via `gss_wrap`) available. `gss_wrap` will provide message encapsulation, data-origin authentication and integrity services only.

GSS_C_INTEG_FLAG

- True - Integrity service may be invoked by calling either `gss_get_mic` or `gss_wrap` routines.
- False - Per-message integrity service unavailable.

GSS_C_ANON_FLAG

- True - The initiator does not wish to be authenticated; the `src_name` parameter (if requested) contains an anonymous internal name.
- False - The initiator has been authenticated normally.

GSS_C_PROT_READY_FLAG

- True - Protection services (as specified by the states of the `GSS_C_CONF_FLAG` and `GSS_C_INTEG_FLAG`) are available if the accompanying major status return value is either `GSS_S_COMPLETE` or `GSS_S_CONTINUE_NEEDED`.
- False - Protection services (as specified by the states of the `GSS_C_CONF_FLAG` and `GSS_C_INTEG_FLAG`) are available only if the accompanying major status return value is `GSS_S_COMPLETE`.

GSS_C_TRANS_FLAG

- True - The resultant security context may be transferred to other processes via a call to `gss_export_sec_context()`.
- False - The security context is not transferable.

All other bits should be set to zero.

Return value:

GSS_S_CONTINUE_NEEDED: Indicates that a token from the peer application is required to complete the context, and that `gss_accept_sec_context` must be called again with that token.

GSS_S_DEFECTIVE_TOKEN: Indicates that consistency checks performed on the `input_token` failed.

GSS_S_DEFECTIVE_CREDENTIAL: Indicates that consistency checks performed on the credential failed.

GSS_S_NO_CRED: The supplied credentials were not valid for context acceptance, or the credential handle did not reference any credentials.

GSS_S_CREDENTIALS_EXPIRED: The referenced credentials have expired.

GSS_S_BAD_BINDINGS: The `input_token` contains different channel bindings to those specified via the `input_chan_bindings` parameter.

GSS_S_NO_CONTEXT: Indicates that the supplied context handle did not refer to a valid context.

GSS_S_BAD_SIG: The `input_token` contains an invalid MIC.

GSS_S_OLD_TOKEN: The `input_token` was too old. This is a fatal error during context establishment.

GSS_S_DUPLICATE_TOKEN: The `input_token` is valid, but is a duplicate of a token already processed. This is a fatal error during context establishment.

GSS_S_BAD_MECH: The received token specified a mechanism that is not supported by the implementation or the provided credential.

`OM_uint32 gss_delete_sec_context (OM_uint32 * minor_status, [Function]
 gss_ctx_id_t * context_handle, gss_buffer_t output_token)`

minor_status: (Integer, modify) Mechanism specific status code.

context_handle: (`gss_ctx_id_t`, modify) Context handle identifying context to delete. After deleting the context, the GSS-API will set this context handle to `GSS_C_NO_CONTEXT`.

output_token: (buffer, opaque, modify, optional) Token to be sent to remote application to instruct it to also delete the context. It is recommended that applications specify `GSS_C_NO_BUFFER` for this parameter, requesting local deletion only. If a buffer parameter is provided by the application, the mechanism may return a token in it; mechanisms that implement only local deletion should set the length field of this token to zero to indicate to the application that no token is to be sent to the peer.

Delete a security context. `gss_delete_sec_context` will delete the local data structures associated with the specified security context, and may generate an `output_token`, which when passed to the peer `gss_process_context_token` will instruct it to do likewise. If no token is required by the mechanism, the GSS-API should set the length field of the `output_token` (if provided) to zero. No further security services may be obtained using the context specified by `context_handle`.

In addition to deleting established security contexts, `gss_delete_sec_context` must also be able to delete "half-built" security contexts resulting from an incomplete sequence of `gss_init_sec_context()/gss_accept_sec_context()` calls.

The `output_token` parameter is retained for compatibility with version 1 of the GSS-API. It is recommended that both peer applications invoke `gss_delete_sec_context` passing the value `GSS_C_NO_BUFFER` for the `output_token` parameter, indicating that no token is required, and that `gss_delete_sec_context` should simply delete local context data structures. If the application does pass a valid buffer to `gss_delete_sec_context`, mechanisms are encouraged to return a zero-length token, indicating that no peer action is necessary, and that no token should be transferred by the application.

Return value:

GSS_S_COMPLETE: Successful completion.

GSS_S_NO_CONTEXT: No valid context was supplied.

`OM_uint32 gss_process_context_token (OM_uint32 * [Function]
minor_status, const gss_ctx_id_t context_handle, const gss_buffer_t token_buffer)`

minor_status: (Integer, modify) Implementation specific status code.

context_handle: (*gss_ctx_id_t*, read) Context handle of context on which token is to be processed

token_buffer: (buffer, opaque, read) Token to process.

Provides a way to pass an asynchronous token to the security service. Most context-level tokens are emitted and processed synchronously by `gss_init_sec_context` and `gss_accept_sec_context`, and the application is informed as to whether further tokens are expected by the `GSS_C_CONTINUE_NEEDED` major status bit. Occasionally, a mechanism may need to emit a context-level token at a point when the peer entity is not expecting a token. For example, the initiator's final call to `gss_init_sec_context` may emit a token and return a status of `GSS_S_COMPLETE`, but the acceptor's call to `gss_accept_sec_context` may fail. The acceptor's mechanism may wish to send a token containing an error indication to the initiator, but the initiator is not expecting a token at this point, believing that the context is fully established. `Gss_process_context_token` provides a way to pass such a token to the mechanism at any time.

Return value:

`GSS_S_COMPLETE`: Successful completion.

`GSS_S_DEFECTIVE_TOKEN`: Indicates that consistency checks performed on the token failed.

`GSS_S_NO_CONTEXT`: The `context_handle` did not refer to a valid context.

`OM_uint32 gss_context_time (OM_uint32 * minor_status, const [Function]
gss_ctx_id_t context_handle, OM_uint32 * time_rec)`

minor_status: (Integer, modify) Implementation specific status code.

context_handle: (*gss_ctx_id_t*, read) Identifies the context to be interrogated.

time_rec: (Integer, modify) Number of seconds that the context will remain valid. If the context has already expired, zero will be returned.

Determines the number of seconds for which the specified context will remain valid.

Return value:

`GSS_S_COMPLETE`: Successful completion.

`GSS_S_CONTEXT_EXPIRED`: The context has already expired.

`GSS_S_NO_CONTEXT`: The `context_handle` parameter did not identify a valid context

`OM_uint32 gss_inquire_context (OM_uint32 * minor_status, const [Function]
gss_ctx_id_t context_handle, gss_name_t * src_name, gss_name_t * targ_name, OM_uint32 * lifetime_rec, gss_OID * mech_type, OM_uint32 * ctx_flags, int * locally_initiated, int * open)`

minor_status: (Integer, modify) Mechanism specific status code.

context_handle: (*gss_ctx_id_t*, read) A handle that refers to the security context.

src_name: (gss_name_t, modify, optional) The name of the context initiator. If the context was established using anonymous authentication, and if the application invoking `gss_inquire_context` is the context acceptor, an anonymous name will be returned. Storage associated with this name must be freed by the application after use with a call to `gss_release_name()`. Specify NULL if not required.

targ_name: (gss_name_t, modify, optional) The name of the context acceptor. Storage associated with this name must be freed by the application after use with a call to `gss_release_name()`. If the context acceptor did not authenticate itself, and if the initiator did not specify a target name in its call to `gss_init_sec_context()`, the value `GSS_C_NO_NAME` will be returned. Specify NULL if not required.

lifetime_rec: (Integer, modify, optional) The number of seconds for which the context will remain valid. If the context has expired, this parameter will be set to zero. If the implementation does not support context expiration, the value `GSS_C_INDEFINITE` will be returned. Specify NULL if not required.

mech_type: (gss_OID, modify, optional) The security mechanism providing the context. The returned OID will be a pointer to static storage that should be treated as read-only by the application; in particular the application should not attempt to free it. Specify NULL if not required.

ctx_flags: (bit-mask, modify, optional) Contains various independent flags, each of which indicates that the context supports (or is expected to support, if `ctx_open` is false) a specific service option. If not needed, specify NULL. Symbolic names are provided for each flag, and the symbolic names corresponding to the required flags should be logically-ANDed with the `ret_flags` value to test whether a given option is supported by the context. See below for the flags.

locally_initiated: (Boolean, modify) Non-zero if the invoking application is the context initiator. Specify NULL if not required.

open: (Boolean, modify) Non-zero if the context is fully established; Zero if a context-establishment token is expected from the peer application. Specify NULL if not required.

Obtains information about a security context. The caller must already have obtained a handle that refers to the context, although the context need not be fully established.

The `ctx_flags` values:

GSS_C_DELEG_FLAG

- True - Credentials were delegated from the initiator to the acceptor.
- False - No credentials were delegated.

GSS_C_MUTUAL_FLAG

- True - The acceptor was authenticated to the initiator.
- False - The acceptor did not authenticate itself.

GSS_C_REPLAY_FLAG

- True - replay of protected messages will be detected.
- False - replayed messages will not be detected.

GSS_C_SEQUENCE_FLAG

- True - out-of-sequence protected messages will be detected.
- False - out-of-sequence messages will not be detected.

GSS_C_CONF_FLAG

- True - Confidentiality service may be invoked by calling `gss_wrap` routine.
- False - No confidentiality service (via `gss_wrap`) available. `gss_wrap` will provide message encapsulation, data-origin authentication and integrity services only.

GSS_C_INTEG_FLAG

- True - Integrity service may be invoked by calling either `gss_get_mic` or `gss_wrap` routines.
- False - Per-message integrity service unavailable.

GSS_C_ANON_FLAG

- True - The initiator's identity will not be revealed to the acceptor. The `src_name` parameter (if requested) contains an anonymous internal name.
- False - The initiator has been authenticated normally.

GSS_C_PROT_READY_FLAG

- True - Protection services (as specified by the states of the `GSS_C_CONF_FLAG` and `GSS_C_INTEG_FLAG`) are available for use.
- False - Protection services (as specified by the states of the `GSS_C_CONF_FLAG` and `GSS_C_INTEG_FLAG`) are available only if the context is fully established (i.e. if the open parameter is non-zero).

GSS_C_TRANS_FLAG

- True - The resultant security context may be transferred to other processes via a call to `gss_export_sec_context()`.
- False - The security context is not transferable.

Return value:

`GSS_S_COMPLETE`: Successful completion.

`GSS_S_NO_CONTEXT`: The referenced context could not be accessed.

```
OM_uint32 gss_wrap_size_limit (OM_uint32 *minor_status, const [Function]
    gss_ctx_id_t context_handle, int conf_req_flag, gss_qop_t qop_req,
    OM_uint32 req_output_size, OM_uint32 *max_input_size)
```

minor_status: (Integer, modify) Mechanism specific status code.

context_handle: (`gss_ctx_id_t`, read) A handle that refers to the security over which the messages will be sent.

conf_req_flag: (Boolean, read) Indicates whether `gss_wrap` will be asked to apply confidentiality protection in addition to integrity protection. See the routine description for `gss_wrap` for more details.

qop_req: (*gss_qop_t*, read) Indicates the level of protection that *gss_wrap* will be asked to provide. See the routine description for *gss_wrap* for more details.

req_output_size: (Integer, read) The desired maximum size for tokens emitted by *gss_wrap*.

max_input_size: (Integer, modify) The maximum input message size that may be presented to *gss_wrap* in order to guarantee that the emitted token shall be no larger than *req_output_size* bytes.

Allows an application to determine the maximum message size that, if presented to *gss_wrap* with the same *conf_req_flag* and *qop_req* parameters, will result in an output token containing no more than *req_output_size* bytes.

This call is intended for use by applications that communicate over protocols that impose a maximum message size. It enables the application to fragment messages prior to applying protection.

GSS-API implementations are recommended but not required to detect invalid QOP values when *gss_wrap_size_limit()* is called. This routine guarantees only a maximum message size, not the availability of specific QOP values for message protection.

Successful completion of this call does not guarantee that *gss_wrap* will be able to protect a message of length *max_input_size* bytes, since this ability may depend on the availability of system resources at the time that *gss_wrap* is called. However, if the implementation itself imposes an upper limit on the length of messages that may be processed by *gss_wrap*, the implementation should not return a value via *max_input_bytes* that is greater than this length.

Return value:

GSS_S_COMPLETE: Successful completion.

GSS_S_NO_CONTEXT: The referenced context could not be accessed.

GSS_S_CONTEXT_EXPIRED: The context has expired.

GSS_S_BAD_QOP: The specified QOP is not supported by the mechanism.

`OM_uint32 gss_export_sec_context (OM_uint32 *minor_status, [Function]
gss_ctx_id_t *context_handle, gss_buffer_t interprocess_token)`

minor_status: (Integer, modify) Mechanism specific status code.

context_handle: (*gss_ctx_id_t*, modify) Context handle identifying the context to transfer.

interprocess_token: (buffer, opaque, modify) Token to be transferred to target process. Storage associated with this token must be freed by the application after use with a call to *gss_release_buffer()*.

Provided to support the sharing of work between multiple processes. This routine will typically be used by the context-acceptor, in an application where a single process receives incoming connection requests and accepts security contexts over them, then passes the established context to one or more other processes for message exchange. *gss_export_sec_context()* deactivates the security context for the calling process and creates an interprocess token which, when passed to *gss_import_sec_context* in another process, will re-activate the context in the second process. Only a single instantiation

of a given context may be active at any one time; a subsequent attempt by a context exporter to access the exported security context will fail.

The implementation may constrain the set of processes by which the interprocess token may be imported, either as a function of local security policy, or as a result of implementation decisions. For example, some implementations may constrain contexts to be passed only between processes that run under the same account, or which are part of the same process group.

The interprocess token may contain security-sensitive information (for example cryptographic keys). While mechanisms are encouraged to either avoid placing such sensitive information within interprocess tokens, or to encrypt the token before returning it to the application, in a typical object-library GSS-API implementation this may not be possible. Thus the application must take care to protect the interprocess token, and ensure that any process to which the token is transferred is trustworthy.

If creation of the interprocess token is successful, the implementation shall deallocate all process-wide resources associated with the security context, and set the `context_handle` to `GSS_C_NO_CONTEXT`. In the event of an error that makes it impossible to complete the export of the security context, the implementation must not return an interprocess token, and should strive to leave the security context referenced by the `context_handle` parameter untouched. If this is impossible, it is permissible for the implementation to delete the security context, providing it also sets the `context_handle` parameter to `GSS_C_NO_CONTEXT`.

Return value:

`GSS_S_COMPLETE`: Successful completion.

`GSS_S_CONTEXT_EXPIRED`: The context has expired.

`GSS_S_NO_CONTEXT`: The context was invalid.

`GSS_S_UNAVAILABLE`: The operation is not supported.

`OM_uint32 gss_import_sec_context (OM_uint32 * minor_status, [Function]
 const gss_buffer_t interprocess_token, gss_ctx_id_t * context_handle)`

minor_status: (Integer, modify) Mechanism specific status code.

interprocess_token: (buffer, opaque, modify) Token received from exporting process

context_handle: (*gss_ctx_id_t*, modify) Context handle of newly reactivated context. Resources associated with this context handle must be released by the application after use with a call to `gss_delete_sec_context()`.

Allows a process to import a security context established by another process. A given interprocess token may be imported only once. See `gss_export_sec_context`.

Return value:

`GSS_S_COMPLETE`: Successful completion.

`GSS_S_NO_CONTEXT`: The token did not contain a valid context reference.

`GSS_S_DEFECTIVE_TOKEN`: The token was invalid.

`GSS_S_UNAVAILABLE`: The operation is unavailable.

`GSS_S_UNAUTHORIZED`: Local policy prevents the import of this context by the current process.

3.7 Per-Message Routines

GSS-API Per-message Routines

Routine	Function
-----	-----
<code>gss_get_mic</code>	Calculate a cryptographic message integrity code (MIC) for a message; integrity service.
<code>gss_verify_mic</code>	Check a MIC against a message; verify integrity of a received message.
<code>gss_wrap</code>	Attach a MIC to a message, and optionally encrypt the message content. confidentiality service
<code>gss_unwrap</code>	Verify a message with attached MIC, and decrypt message content if necessary.

`OM_uint32 gss_get_mic (OM_uint32 * minor_status, const [Function]
gss_ctx_id_t context_handle, gss_qop_t qop_req, const gss_buffer_t message_buffer, gss_buffer_t message_token)`

minor_status: (Integer, modify) Mechanism specific status code.

context_handle: (*gss_ctx_id_t*, read) Identifies the context on which the message will be sent.

qop_req: (*gss_qop_t*, read, optional) Specifies requested quality of protection. Callers are encouraged, on portability grounds, to accept the default quality of protection offered by the chosen mechanism, which may be requested by specifying `GSS_C_QOP_DEFAULT` for this parameter. If an unsupported protection strength is requested, `gss_get_mic` will return a `major_status` of `GSS_S_BAD_QOP`.

message_buffer: (buffer, opaque, read) Message to be protected.

message_token: (buffer, opaque, modify) Buffer to receive token. The application must free storage associated with this buffer after use with a call to `gss_release_buffer()`.

Generates a cryptographic MIC for the supplied message, and places the MIC in a token for transfer to the peer application. The `qop_req` parameter allows a choice between several cryptographic algorithms, if supported by the chosen mechanism.

Since some application-level protocols may wish to use tokens emitted by `gss_wrap()` to provide "secure framing", implementations must support derivation of MICs from zero-length messages.

Return value:

`GSS_S_COMPLETE`: Successful completion.

`GSS_S_CONTEXT_EXPIRED`: The context has already expired.

`GSS_S_NO_CONTEXT`: The `context_handle` parameter did not identify a valid context.

`GSS_S_BAD_QOP`: The specified QOP is not supported by the mechanism.

```
OM_uint32 gss_verify_mic (OM_uint32 * minor_status, const [Function]
    gss_ctx_id_t context_handle, const gss_buffer_t message_buffer, const
    gss_buffer_t token_buffer, gss_qop_t * qop_state)
```

minor_status: (Integer, modify) Mechanism specific status code.

context_handle: (gss_ctx_id_t, read) Identifies the context on which the message arrived.

message_buffer: (buffer, opaque, read) Message to be verified.

token_buffer: (buffer, opaque, read) Token associated with message.

qop_state: (gss_qop_t, modify, optional) Quality of protection gained from MIC Specify NULL if not required.

Verifies that a cryptographic MIC, contained in the token parameter, fits the supplied message. The *qop_state* parameter allows a message recipient to determine the strength of protection that was applied to the message.

Since some application-level protocols may wish to use tokens emitted by `gss_wrap()` to provide "secure framing", implementations must support the calculation and verification of MICs over zero-length messages.

Return value:

GSS_S_COMPLETE: Successful completion.

GSS_S_DEFECTIVE_TOKEN: The token failed consistency checks.

GSS_S_BAD_SIG: The MIC was incorrect.

GSS_S_DUPLICATE_TOKEN: The token was valid, and contained a correct MIC for the message, but it had already been processed.

GSS_S_OLD_TOKEN: The token was valid, and contained a correct MIC for the message, but it is too old to check for duplication.

GSS_S_UNSEQ_TOKEN: The token was valid, and contained a correct MIC for the message, but has been verified out of sequence; a later token has already been received.

GSS_S_GAP_TOKEN: The token was valid, and contained a correct MIC for the message, but has been verified out of sequence; an earlier expected token has not yet been received.

GSS_S_CONTEXT_EXPIRED: The context has already expired.

GSS_S_NO_CONTEXT: The *context_handle* parameter did not identify a valid context.

```
OM_uint32 gss_wrap (OM_uint32 * minor_status, const gss_ctx_id_t [Function]
    context_handle, int conf_req_flag, gss_qop_t qop_req, const gss_buffer_t
    input_message_buffer, int * conf_state, gss_buffer_t
    output_message_buffer)
```

minor_status: (Integer, modify) Mechanism specific status code.

context_handle: (gss_ctx_id_t, read) Identifies the context on which the message will be sent.

conf_req_flag: (boolean, read) Non-zero - Both confidentiality and integrity services are requested. Zero - Only integrity service is requested.

qop_req: (gss_qop_t, read, optional) Specifies required quality of protection. A mechanism-specific default may be requested by setting *qop_req* to

GSS_C_QOP_DEFAULT. If an unsupported protection strength is requested, `gss_wrap` will return a `major_status` of `GSS_S_BAD_QOP`.

input_message_buffer: (buffer, opaque, read) Message to be protected.

conf_state: (boolean, modify, optional) Non-zero - Confidentiality, data origin authentication and integrity services have been applied. Zero - Integrity and data origin services only has been applied. Specify NULL if not required.

output_message_buffer: (buffer, opaque, modify) Buffer to receive protected message. Storage associated with this message must be freed by the application after use with a call to `gss_release_buffer()`.

Attaches a cryptographic MIC and optionally encrypts the specified `input_message`. The `output_message` contains both the MIC and the message. The `qop_req` parameter allows a choice between several cryptographic algorithms, if supported by the chosen mechanism.

Since some application-level protocols may wish to use tokens emitted by `gss_wrap()` to provide "secure framing", implementations must support the wrapping of zero-length messages.

Return value:

GSS_S_COMPLETE: Successful completion.

GSS_S_CONTEXT_EXPIRED: The context has already expired.

GSS_S_NO_CONTEXT: The `context_handle` parameter did not identify a valid context.

GSS_S_BAD_QOP: The specified QOP is not supported by the mechanism.

```
OM_uint32 gss_unwrap (OM_uint32 *minor_status, const [Function]
                    gss_ctx_id_t context_handle, const gss_buffer_t input_message_buffer,
                    gss_buffer_t output_message_buffer, int * conf_state, gss_qop_t *
                    qop_state)
```

minor_status: (Integer, modify) Mechanism specific status code.

context_handle: (`gss_ctx_id_t`, read) Identifies the context on which the message arrived.

input_message_buffer: (buffer, opaque, read) Protected message.

output_message_buffer: (buffer, opaque, modify) Buffer to receive unwrapped message. Storage associated with this buffer must be freed by the application after use with a call to `gss_release_buffer()`.

conf_state: (boolean, modify, optional) Non-zero - Confidentiality and integrity protection were used. Zero - Integrity service only was used. Specify NULL if not required.

qop_state: (`gss_qop_t`, modify, optional) Quality of protection provided. Specify NULL if not required.

Converts a message previously protected by `gss_wrap` back to a usable form, verifying the embedded MIC. The `conf_state` parameter indicates whether the message was encrypted; the `qop_state` parameter indicates the strength of protection that was used to provide the confidentiality and integrity services.

Since some application-level protocols may wish to use tokens emitted by `gss_wrap()` to provide "secure framing", implementations must support the wrapping and un-wrapping of zero-length messages.

Return value:

`GSS_S_COMPLETE`: Successful completion.

`GSS_S_DEFECTIVE_TOKEN`: The token failed consistency checks.

`GSS_S_BAD_SIG`: The MIC was incorrect.

`GSS_S_DUPLICATE_TOKEN`: The token was valid, and contained a correct MIC for the message, but it had already been processed.

`GSS_S_OLD_TOKEN`: The token was valid, and contained a correct MIC for the message, but it is too old to check for duplication.

`GSS_S_UNSEQ_TOKEN`: The token was valid, and contained a correct MIC for the message, but has been verified out of sequence; a later token has already been received.

`GSS_S_GAP_TOKEN`: The token was valid, and contained a correct MIC for the message, but has been verified out of sequence; an earlier expected token has not yet been received.

`GSS_S_CONTEXT_EXPIRED`: The context has already expired.

`GSS_S_NO_CONTEXT`: The `context_handle` parameter did not identify a valid context.

3.8 Name Manipulation

GSS-API Name manipulation Routines

Routine	Function
-----	-----
<code>gss_import_name</code>	Convert a contiguous string name to internal-form.
<code>gss_display_name</code>	Convert internal-form name to text.
<code>gss_compare_name</code>	Compare two internal-form names.
<code>gss_release_name</code>	Discard an internal-form name.
<code>gss_inquire_names_for_mech</code>	List the name-types supported by the specified mechanism.
<code>gss_inquire_mechs_for_name</code>	List mechanisms that support the specified name-type.
<code>gss_canonicalize_name</code>	Convert an internal name to an MN.
<code>gss_export_name</code>	Convert an MN to export form.
<code>gss_duplicate_name</code>	Create a copy of an internal name.

```
OM_uint32 gss_import_name (OM_uint32 * minor_status, const [Function]
    gss_buffer_t input_name_buffer, const gss_OID input_name_type,
    gss_name_t * output_name)
```

minor_status: (Integer, modify) Mechanism specific status code.

input_name_buffer: (buffer, octet-string, read) Buffer containing contiguous string name to convert.

input_name_type: (Object ID, read, optional) Object ID specifying type of printable name. Applications may specify either GSS_C_NO_OID to use a mechanism-specific default printable syntax, or an OID recognized by the GSS-API implementation to name a specific namespace.

output_name: (*gss_name_t*, modify) Returned name in internal form. Storage associated with this name must be freed by the application after use with a call to *gss_release_name()*.

Convert a contiguous string name to internal form. In general, the internal name returned (via the *@output_name* parameter) will not be an MN; the exception to this is if the *@input_name_type* indicates that the contiguous string provided via the *@input_name_buffer* parameter is of type GSS_C_NT_EXPORT_NAME, in which case the returned internal name will be an MN for the mechanism that exported the name.

Return value:

GSS_S_COMPLETE: Successful completion.

GSS_S_BAD_NAME_TYPE: The *input_name_type* was unrecognized.

GSS_S_BAD_NAME: The *input_name* parameter could not be interpreted as a name of the specified type.

GSS_S_BAD_MECH: The input name-type was GSS_C_NT_EXPORT_NAME, but the mechanism contained within the input-name is not supported.

```
OM_uint32 gss_display_name (OM_uint32 *minor_status, const [Function]
    gss_name_t input_name, gss_buffer_t output_name_buffer, gss_OID *
    output_name_type)
```

minor_status: (Integer, modify) Mechanism specific status code.

input_name: (*gss_name_t*, read) Name to be displayed.

output_name_buffer: (buffer, character-string, modify) Buffer to receive textual name string. The application must free storage associated with this name after use with a call to *gss_release_buffer()*.

output_name_type: (Object ID, modify, optional) The type of the returned name. The returned *gss_OID* will be a pointer into static storage, and should be treated as read-only by the caller (in particular, the application should not attempt to free it). Specify NULL if not required.

Allows an application to obtain a textual representation of an opaque internal-form name for display purposes. The syntax of a printable name is defined by the GSS-API implementation.

If *input_name* denotes an anonymous principal, the implementation should return the *gss_OID* value GSS_C_NT_ANONYMOUS as the *output_name_type*, and a textual name that is syntactically distinct from all valid supported printable names in *output_name_buffer*.

If *input_name* was created by a call to *gss_import_name*, specifying GSS_C_NO_OID as the name-type, implementations that employ lazy conversion between name types may return GSS_C_NO_OID via the *output_name_type* parameter.

Return value:

GSS_S_COMPLETE: Successful completion.

GSS_S_BAD_NAME: *@input_name* was ill-formed.

`OM_uint32 gss_compare_name (OM_uint32 *minor_status, const gss_name_t name1, const gss_name_t name2, int *name_equal)` [Function]

minor_status: (Integer, modify) Mechanism specific status code.

name1: (gss_name_t, read) Internal-form name.

name2: (gss_name_t, read) Internal-form name.

name_equal: (boolean, modify) Non-zero - names refer to same entity. Zero - names refer to different entities (strictly, the names are not known to refer to the same identity).

Allows an application to compare two internal-form names to determine whether they refer to the same entity.

If either name presented to `gss_compare_name` denotes an anonymous principal, the routines should indicate that the two names do not refer to the same identity.

Return value:

`GSS_S_COMPLETE`: Successful completion.

`GSS_S_BAD_NAME_TYPE`: The two names were of incomparable types.

`GSS_S_BAD_NAME`: One or both of *name1* or *name2* was ill-formed.

`OM_uint32 gss_release_name (OM_uint32 *minor_status, gss_name_t *name)` [Function]

minor_status: (Integer, modify) Mechanism specific status code.

name: (gss_name_t, modify) The name to be deleted.

Free GSSAPI-allocated storage associated with an internal-form name. The name is set to `GSS_C_NO_NAME` on successful completion of this call.

Return value:

`GSS_S_COMPLETE`: Successful completion.

`GSS_S_BAD_NAME`: The name parameter did not contain a valid name.

`OM_uint32 gss_inquire_names_for_mech (OM_uint32 *minor_status, const gss_OID mechanism, gss_OID_set *name_types)` [Function]

minor_status: (Integer, modify) Mechanism specific status code.

mechanism: (gss_OID, read) The mechanism to be interrogated.

name_types: (gss_OID_set, modify) Set of name-types supported by the specified mechanism. The returned OID set must be freed by the application after use with a call to `gss_release_oid_set()`.

Returns the set of nametypes supported by the specified mechanism.

Return value:

`GSS_S_COMPLETE`: Successful completion.

`OM_uint32 gss_inquire_mechs_for_name (OM_uint32 *minor_status, const gss_name_t input_name, gss_OID_set *mech_types)` [Function]

minor_status: (Integer, modify) Mechanism specific status code.

input_name: (gss_name_t, read) The name to which the inquiry relates.

mech_types: (gss_OID_set, modify) Set of mechanisms that may support the specified name. The returned OID set must be freed by the caller after use with a call to `gss_release_oid_set()`.

Returns the set of mechanisms supported by the GSS-API implementation that may be able to process the specified name.

Each mechanism returned will recognize at least one element within the name. It is permissible for this routine to be implemented within a mechanism-independent GSS-API layer, using the type information contained within the presented name, and based on registration information provided by individual mechanism implementations. This means that the returned *mech_types* set may indicate that a particular mechanism will understand the name when in fact it would refuse to accept the name as input to `gss_canonicalize_name`, `gss_init_sec_context`, `gss_acquire_cred` or `gss_add_cred` (due to some property of the specific name, as opposed to the name type). Thus this routine should be used only as a prefilter for a call to a subsequent mechanism-specific routine.

Return value:

GSS_S_COMPLETE: Successful completion.

GSS_S_BAD_NAME: The *input_name* parameter was ill-formed.

GSS_S_BAD_NAME_TYPE: The *input_name* parameter contained an invalid or unsupported type of name.

```
OM_uint32 gss_canonicalize_name (OM_uint32 *minor_status,      [Function]
                                const gss_name_t input_name, const gss_OID mech_type, gss_name_t *
                                output_name)
```

minor_status: (Integer, modify) Mechanism specific status code.

input_name: (gss_name_t, read) The name for which a canonical form is desired.

mech_type: (Object ID, read) The authentication mechanism for which the canonical form of the name is desired. The desired mechanism must be specified explicitly; no default is provided.

output_name: (gss_name_t, modify) The resultant canonical name. Storage associated with this name must be freed by the application after use with a call to `gss_release_name()`.

Generate a canonical mechanism name (MN) from an arbitrary internal name. The mechanism name is the name that would be returned to a context acceptor on successful authentication of a context where the initiator used the *input_name* in a successful call to `gss_acquire_cred`, specifying an OID set containing @*mech_type* as its only member, followed by a call to `gss_init_sec_context()`, specifying @*mech_type* as the authentication mechanism.

Return value:

GSS_S_COMPLETE: Successful completion.

```
OM_uint32 gss_export_name (OM_uint32 *minor_status, const      [Function]
                           gss_name_t input_name, gss_buffer_t exported_name)
```

minor_status: (Integer, modify) Mechanism specific status code.

input_name: (gss_name_t, read) The MN to be exported.

exported_name: (gss_buffer_t, octet-string, modify) The canonical contiguous string form of @input_name. Storage associated with this string must be freed by the application after use with gss_release_buffer().

To produce a canonical contiguous string representation of a mechanism name (MN), suitable for direct comparison (e.g. with memcmp) for use in authorization functions (e.g. matching entries in an access-control list). The @input_name parameter must specify a valid MN (i.e. an internal name generated by gss_accept_sec_context() or by gss_canonicalize_name()).

Return value:

GSS_S_COMPLETE: Successful completion.

GSS_S_NAME_NOT_MN: The provided internal name was not a mechanism name.

GSS_S_BAD_NAME: The provided internal name was ill-formed.

GSS_S_BAD_NAME_TYPE: The internal name was of a type not supported by the GSS-API implementation.

OM_uint32 gss_duplicate_name (OM_uint32 * minor_status, const [Function]
gss_name_t src_name, gss_name_t * dest_name)

minor_status: (Integer, modify) Mechanism specific status code.

src_name: (gss_name_t, read) Internal name to be duplicated.

dest_name: (gss_name_t, modify) The resultant copy of @src_name. Storage associated with this name must be freed by the application after use with a call to gss_release_name().

Create an exact duplicate of the existing internal name @src_name. The new @dest_name will be independent of src_name (i.e. @src_name and @dest_name must both be released, and the release of one shall not affect the validity of the other).

Return value:

GSS_S_COMPLETE: Successful completion.

GSS_S_BAD_NAME: The src_name parameter was ill-formed.

3.9 Miscellaneous Routines

GSS-API Miscellaneous Routines

Routine	Function
-----	-----
gss_add_oid_set_member	Add an object identifier to a set.
gss_display_status	Convert a GSS-API status code to text.
gss_indicate_mechs	Determine available underlying authentication mechanisms.
gss_release_buffer	Discard a buffer.
gss_release_oid_set	Discard a set of object identifiers.
gss_create_empty_oid_set	Create a set containing no

object identifiers.
 gss_test_oid_set_member Determines whether an object identifier is a member of a set.

OM_uint32 gss_add_oid_set_member (OM_uint32 *minor_status, [Function]
 const gss_OID member_oid, gss_OID_set *oid_set)
minor_status: (integer, modify) Mechanism specific status code.
member_oid: (Object ID, read) The object identifier to copied into the set.
oid_set: (Set of Object ID, modify) The set in which the object identifier should be inserted.

Add an Object Identifier to an Object Identifier set. This routine is intended for use in conjunction with gss_create_empty_oid_set when constructing a set of mechanism OIDs for input to gss_acquire_cred. The oid_set parameter must refer to an OID-set that was created by GSS-API (e.g. a set returned by gss_create_empty_oid_set()). GSS-API creates a copy of the member_oid and inserts this copy into the set, expanding the storage allocated to the OID-set's elements array if necessary. The routine may add the new member OID anywhere within the elements array, and implementations should verify that the new member_oid is not already contained within the elements array; if the member_oid is already present, the oid_set should remain unchanged.

Return value:

GSS_S_COMPLETE: Successful completion.

OM_uint32 gss_display_status (OM_uint32 *minor_status, [Function]
 OM_uint32 status_value, int status_type, const gss_OID mech_type,
 OM_uint32 *message_context, gss_buffer_t status_string)
minor_status: (integer, modify) Mechanism specific status code.
status_value: (Integer, read) Status value to be converted.
status_type: (Integer, read) GSS_C_GSS_CODE - status_value is a GSS status code.
 GSS_C_MECH_CODE - status_value is a mechanism status code.
mech_type: (Object ID, read, optional) Underlying mechanism (used to interpret a minor status value). Supply GSS_C_NO_OID to obtain the system default.
message_context: (Integer, read/modify) Should be initialized to zero by the application prior to the first call. On return from gss_display_status(), a non-zero status_value parameter indicates that additional messages may be extracted from the status code via subsequent calls to gss_display_status(), passing the same status_value, status_type, mech_type, and message_context parameters.
status_string: (buffer, character string, modify) Textual interpretation of the status_value. Storage associated with this parameter must be freed by the application after use with a call to gss_release_buffer().

Allows an application to obtain a textual representation of a GSS-API status code, for display to the user or for logging purposes. Since some status values may indicate multiple conditions, applications may need to call gss_display_status multiple times, each call generating a single text string. The message_context parameter is used by gss_display_status to store state information about which error messages have already been extracted from a given status_value; message_context must be initialized to 0

by the application prior to the first call, and `gss_display_status` will return a non-zero value in this parameter if there are further messages to extract.

The `message_context` parameter contains all state information required by `gss_display_status` in order to extract further messages from the `status_value`; even when a non-zero value is returned in this parameter, the application is not required to call `gss_display_status` again unless subsequent messages are desired. The following code extracts all messages from a given status code and prints them to `stderr`:

```

OM_uint32 message_context;
OM_uint32 status_code;
OM_uint32 maj_status;
OM_uint32 min_status;
gss_buffer_desc status_string;

    ...

message_context = 0;

do {
    maj_status = gss_display_status (
                &min_status,
                status_code,
                GSS_C_GSS_CODE,
                GSS_C_NO_OID,
                &message_context,
                &status_string)

    fprintf(stderr,
            "%.*s\n",
            (int)status_string.length,
            (char *)status_string.value);

    gss_release_buffer(&min_status, &status_string);

} while (message_context != 0);

```

Return value:

GSS_S_COMPLETE: Successful completion.

GSS_S_BAD_MECH: Indicates that translation in accordance with an unsupported mechanism type was requested.

GSS_S_BAD_STATUS: The status value was not recognized, or the status type was neither `GSS_C_GSS_CODE` nor `GSS_C_MECH_CODE`.

`OM_uint32 gss_indicate_mechs (OM_uint32 * minor_status, [Function]
 gss_OID_set * mech_set)
 minor_status:` (integer, modify) Mechanism specific status code.

mech_set: (set of Object IDs, modify) Set of implementation-supported mechanisms. The returned `gss_OID_set` value will be a dynamically-allocated OID set, that should be released by the caller after use with a call to `gss_release_oid_set()`.

Allows an application to determine which underlying security mechanisms are available.

Return value:

`GSS_S_COMPLETE`: Successful completion.

`OM_uint32 gss_release_buffer (OM_uint32 * minor_status, [Function]
gss_buffer_t buffer)`

minor_status: (integer, modify) Mechanism specific status code.

buffer: (buffer, modify) The storage associated with the buffer will be deleted. The `gss_buffer_desc` object will not be freed, but its length field will be zeroed.

Free storage associated with a buffer. The storage must have been allocated by a GSS-API routine. In addition to freeing the associated storage, the routine will zero the length field in the descriptor to which the buffer parameter refers, and implementations are encouraged to additionally set the pointer field in the descriptor to NULL. Any buffer object returned by a GSS-API routine may be passed to `gss_release_buffer` (even if there is no storage associated with the buffer).

Return value:

`GSS_S_COMPLETE`: Successful completion.

`OM_uint32 gss_release_oid_set (OM_uint32 * minor_status, [Function]
gss_OID_set * set)`

minor_status: (integer, modify) Mechanism specific status code.

set: (Set of Object IDs, modify) The storage associated with the `gss_OID_set` will be deleted.

Free storage associated with a GSSAPI-generated `gss_OID_set` object. The `set` parameter must refer to an OID-set that was returned from a GSS-API routine. `gss_release_oid_set()` will free the storage associated with each individual member OID, the OID set's elements array, and the `gss_OID_set_desc`.

The `gss_OID_set` parameter is set to `GSS_C_NO_OID_SET` on successful completion of this routine.

Return value:

`GSS_S_COMPLETE`: Successful completion.

`OM_uint32 gss_create_empty_oid_set (OM_uint32 * [Function]
minor_status, gss_OID_set * oid_set)`

minor_status: (integer, modify) Mechanism specific status code.

oid_set: (Set of Object IDs, modify) The empty object identifier set. The routine will allocate the `gss_OID_set_desc` object, which the application must free after use with a call to `gss_release_oid_set()`.

Create an object-identifier set containing no object identifiers, to which members may be subsequently added using the `gss_add_oid_set_member()` routine. These routines

are intended to be used to construct sets of mechanism object identifiers, for input to `gss_acquire_cred`.

Return value:

`GSS_S_COMPLETE`: Successful completion.

`OM_uint32 gss_test_oid_set_member (OM_uint32 *minor_status, [Function]
const gss_OID member, const gss_OID_set set, int *present)`

minor_status: (integer, modify) Mechanism specific status code.

member: (Object ID, read) The object identifier whose presence is to be tested.

set: (Set of Object ID, read) The Object Identifier set.

present: (Boolean, modify) Non-zero if the specified OID is a member of the set, zero if not.

Interrogate an Object Identifier set to determine whether a specified Object Identifier is a member. This routine is intended to be used with OID sets returned by `gss_indicate_mechs()`, `gss_acquire_cred()`, and `gss_inquire_cred()`, but will also work with user-generated sets.

Return value:

`GSS_S_COMPLETE`: Successful completion.

4 Extended GSS API

None of the following functions are standard GSS API functions. As such, they are not declared in 'gss/api.h', but rather in 'gss/ext.h' (which is included from 'gss.h'). See [Section 2.1 \[Header\]](#), page 7.

const char * gss_check_version (*const char * req_version*) [Function]
req_version: version string to compare with, or NULL

Check that the the version of the library is at minimum the one given as a string in @req_version.

int gss_oid_equal (*gss_OID first_oid, gss_OID second_oid*) [Function]
first_oid: (Object ID, read) First Object identifier.

second_oid: (Object ID, read) First Object identifier.

Compare two OIDs for equality. The comparison is "deep", i.e., the actual byte sequences of the OIDs are compared instead of just the pointer equality.

OM_uint32 gss_copy_oid (*OM_uint32 * minor_status, const gss_OID src_oid, gss_OID dest_oid*) [Function]

minor_status: (integer, modify) Mechanism specific status code.

src_oid: (Object ID, read) The object identifier to copy.

dest_oid: (Object ID, modify) The resultant copy of @src_oid. Storage associated with this name must be freed by the application, but gss_release_oid() cannot be used generally as it deallocate the the oid structure itself too (use gss_duplicate_oid() if you don't want this problem.)

Make an exact copy of the given OID, that shares no memory areas with the original.

OM_uint32 gss_duplicate_oid (*OM_uint32 * minor_status, const gss_OID src_oid, gss_OID * dest_oid*) [Function]

minor_status: (integer, modify) Mechanism specific status code.

src_oid: (Object ID, read) The object identifier to duplicate.

dest_oid: (Object ID, modify) The resultant copy of @src_oid. Storage associated with this name must be freed by the application, by calling gss_release_oid().

Allocate a new OID and make it an exact copy of the given OID, that shares no memory areas with the original.

int gss_userok (*const gss_name_t name, const char * username*) [Function]
name: (gss_name_t, read) Name to be compared.

username: Zero terminated string with username.

Compare the username against the output from gss_export_name() invoked on @name, after removing the leading OID. This answers the question whether the particular mechanism would authenticate them as the same principal

int gss_encapsulate_token (*gss_buffer_t input_message, gss_OID token_oid, gss_buffer_t output_message*) [Function]

input_message: Message to be encapsulated.

token_oid: OID of mechanism.

input_message: Output buffer with encapsulated message.

Wrap a buffer in the mechanism-independent token format. This is used for the initial token of a GSS-API context establishment sequence. It incorporates an identifier of the mechanism type to be used on that context, and enables tokens to be interpreted unambiguously at GSS-API peers. See further section 3.1 of RFC 2743.

```
int gss_decapsulate_token (gss_buffer_t input_message, gss_OID token_oid, gss_buffer_t output_message) [Function]
```

input_message: Message to decapsulated.

token_oid: Output buffer with mechanism OID used in message.

input_message: Output buffer with encapsulated message.

Unwrap a buffer in the mechanism-independent token format. This is the reverse of `gss_encapsulate_token`. The translation is loss-less, all data is preserved as is.

5 Acknowledgements

This manual borrows text from RFC 2743 and RFC 2744 that describe GSS API formally.

Appendix A Criticism of GSS

The author has doubts whether GSS is the best solution for free software projects looking for a implementation agnostic security framework. We express these doubts in this section, so that the reader can judge for herself if any of the potential problems discussed here are relevant for their project, or if the benefit outweigh the problems. We are aware that some of the opinions are highly subjective, but we offer them in the hope they can serve as anecdotal evidence.

GSS can be criticized on several levels. We start with the actual implementation.

GSS does not appear to be designed by experienced C programmers. While generally this may be a good thing (C is not the best language), but since they defined the API in C, it is unfortunate. The primary evidence of this is the `major_status` and `minor_status` error code solution. It is a complicated way to describe error conditions, but what makes matters worse, the error condition is separated; half of the error condition is in the function return value and the other half is in the first argument to the function, which is always a pointer to an integer. (The pointer is not even allowed to be `NULL`, if the application doesn't care about the minor error code.) This makes the API unreadable, and difficult to use. A better solutions would be to return a struct containing the entire error condition, which can be accessed using macros, although we acknowledge that the C language used at the time GSS was designed may not have allowed this (this may in fact be the reason the awkward solution was chosen). Instead, the return value could have been passed back to callers using a pointer to a struct, accessible using various macros, and the function could have a void prototype. The fact that `minor_status` is placed first in the parameter list increases the pain it is to use the API. Important parameters should be placed first. A better place for `minor_status` (if it must be present at all) would have been last in the prototypes.

Another evidence of the C inexperience are the memory management issues; GSS provides functions to deallocate data stored within, e.g., `gss_buffer_t` but the caller is responsible of deallocating the structure pointed at by the `gss_buffer_t` (i.e., the `gss_buffer_desc`) itself. Memory management issues are error prone, and this division easily leads to memory leaks (or worse). Instead, the API should be the sole owner of all `gss_ctx_id_t`, `gss_cred_id_t`, and `gss_buffer_t` structures: they should be allocated by the library, and deallocated (using the utility functions defined for this purpose) by the library.

TBA: specification is unclear how memory for OIDs are managed. For example, who is responsible for deallocate potentially newly allocated OIDs returned as `actual_mechs` in `gss_acquire_cred`? Further, are OIDs deeply copied into OID sets? In other words, if I add an OID into an OID set, and modify the original OID, will the OID in the OID set be modified too?

Another illustrating example is the sample GSS header file given in the RFC, which contains:

```
/*
 * We have included the xom.h header file.  Verify that OM_uint32
 * is defined correctly.
 */
#if sizeof(gss_uint32) != sizeof(OM_uint32)
#error Incompatible definition of OM_uint32 from xom.h
```

`#endif`

The C pre-processor does not know about the `sizeof` function, so it is treated as an identifier, which maps to 0. Thus, the expression does not check that the size of `OM_uint32` is correct. It checks whether the expression `0 != 0` holds.

TBA: thread issues

TBA: multiple mechanisms in a GSS library

TBA: high-level design criticism.

TBA: no credential forwarding.

TBA: internationalization

TBA: dynamically generated OIDs and memory deallocation issue. I.e., should `gss_import_name` or `gss_duplicate_name` allocate memory and copy the OID provided, or simply copy the pointer? If the former, who would deallocate that memory? If the latter, the application may deallocate or modify the OID, which seem unwanted.

TBA: krb5: no way to access authorization-data

TBA: krb5: firewall/pre-IP: iakerb status?

TBA: krb5: single-DES only

TBA: the API may block, unusable in `select()` based servers. Especially if the servers contacted is decided by the, yet unauthenticated, remote client.

TBA: krb5: no support for `GSS_C_PROT_READY_FLAG`. We support it anyway, though.

TBA: krb5: `gssapi-cfx` differ from rfc 1964 in the reply token in that the latter require presence of sequence numbers whereas the former doesn't.

Finally we note that few free security applications uses GSS, perhaps the only major exception to this are Kerberos 5 implementations. While not substantial evidence, this do suggest that the GSS may not be the simplest solution available to solve actual problems, since otherwise more projects would have chosen to take advantage of the work that went into GSS instead of using another framework (or designing their own solution).

Our conclusion is that free software projects that are looking for a security framework should evaluate carefully whether GSS actually is the best solution before using it. In particular it is recommended to compare GSS with the Simple Authentication and Security Layer (SASL) framework, which in several situations provide the same feature as GSS does. The most compelling argument for SASL over GSS is, as its acronym suggest, Simple, whereas GSS is far from it.

However, that said, for free software projects that wants to support Kerberos 5, we do acknowledge that no other framework provides a more portable and interoperable interface into the Kerberos 5 system. If your project needs to use Kerberos 5 specifically, we do recommend you to use GSS instead of the Kerberos 5 implementation specific APIs.

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Version 1.2, November 2002

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